




Junos® Fundamentals Series

DAY ONE: SECURING THE ROUTING ENGINE ON M, MX, AND T SERIES

A large graphic consisting of overlapping, semi-transparent blue triangles and polygons of various shades, creating a complex, abstract geometric pattern that fills the lower half of the page.

Apply the powerful policy tools of Junos essential to protecting your device and your whole network with expert, step-by-step techniques.

By Douglas Hanks Jr.

DAY ONE: SECURING THE ROUTING ENGINE ON M, MX, AND T SERIES

The routing engine on Junos routers performs many different functions, from processing routing protocol updates, to driving the command-line interface (CLI). Given that the routing engine is critical to the operation of the device and its network, you need to protect the routing engine from unwanted traffic by allowing only essential permitted traffic. Unwanted traffic can come in many different forms: malicious traffic seeking to gain unauthorized access, unintentional routing protocol updates from neighboring devices, or even legitimate traffic that exceeds a given bandwidth limit.

This Day One book shows you how to secure the routing engine step-by-step. Learn how, learn why, then follow along as you build a modular firewall filter and apply it.

“An indispensable resource for anyone who needs to protect their Internet connected routers.”

Matt Hite, Network Engineer, Zynga

IT'S DAY ONE AND YOU HAVE A JOB TO DO, SO LEARN HOW TO:

- Secure the routing engine with a modular framework of firewall policies and policers.
- Understand how firewall filters work and how they are applied to the routing engine.
- Create a modular framework by chaining together firewall filters.
- Describe how firewall logs are managed and view firewall logs in detail.
- Understand how firewall counters work and view firewall counters.
- Write detailed firewall policies to permit specific traffic to the routing engine.
- Build dynamic prefix-lists based off the Junos configuration using apply-path.
- Rate-limit and police certain types of traffic to the routing engine.
- Create filter-specific and term-specific policers.

Juniper Networks Day One books provide just the information you need to know on day one. That's because they are written by subject matter experts who specialize in getting networks up and running. Visit www.juniper.net/dayone to peruse the complete library.

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Day One: Securing the Routing Engine on M, MX, and T Series

By Douglas Hanks Jr.

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Send your suggestions, comments, and critiques by email to dayone@juniper.net.

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What You Need to Know Before Reading this Book

- The reader is expected to have previous hands-on experience working with the Junos operating system and network devices. The majority of this book deals with actual Junos configuration.
- It's beneficial, though not required, for the reader to hold a JNCIA-Junos certification from Juniper Networks. Topics in this book build on the basics found in the JNCIA-Junos material and extend into the intermediate to expert level of configuration.
- An understanding of the services and protocols required to operate your network is necessary. Depending on the function and role of your router, these services and protocols may change. For example, a core router in an enterprise data center may be running OSPF, BFD, and VRRP, whereas an edge router may be running OSPF and BGP.
- A general understanding of Layer 4 protocols such as TCP and UDP is also important. This book focuses heavily on segment fields such as source port, destination port, and flags.

After Reading this Book, You'll Be Able to...

- Secure the routing engine using a modular framework of firewall policies and policers.
- Understand how firewall filters work and how they are applied to the routing engine.
- Create a modular framework by chaining together firewall filters.
- Understand how firewall logs are managed and view them in detail.
- Understand how firewall counters work and view them.
- Write detailed firewall policies to permit specific traffic to the routing engine.
- Build dynamic prefix-lists based on the Junos configuration using the apply-path feature.
- Omit large sections of the Junos configuration using the apply-flags feature.
- Rate-limit and police certain types of traffic to the routing engine.
- Create filter-specific and term-specific policers.

Why Secure the Routing Engine?

The routing engine performs many different functions, from processing routing protocol updates, to driving the command-line interface (CLI). Given that the routing engine is critical for the operation of the router and the network, the routing engine needs to be protected from unwanted traffic and to only allow traffic that it's going to be interested in. Unwanted traffic comes in many different forms: malicious traffic seeking to gain unauthorized access, unintentional routing protocol updates from neighboring devices, or even legitimate traffic that exceeds a given bandwidth limit.

Routers are deployed in a variety of different scenarios and it's common that no two routers are alike, even within the same network. For instance, a router could be deployed in any of the following deployment scenarios:

- Service Provider core, aggregation, and edge.
- Internet transit and peering.
- Enterprise data center and demilitarized zone (DMZ).
- Inter-data center transport.

Each of these deployment scenarios represents different networking requirements in terms of security, protocols, and services, making the creation and maintenance of customized security policies for each router a daunting task.

Even inside a simple enterprise data center there are several distinct components: core, distribution, WAN, and Internet edge. Each of these components will use different networking protocols and services.

Table P.1 Matrix of Components, Protocols, and Services

Component	Protocols	Services
Core	OSPF, BFD, and BGP	RADIUS, SSH, and Web
Distribution	Spanning-tree, BFD, OSPF, and VRRP	RADIUS, SSH, Web, and NTP.
WAN	OSPF, BFP, and BGP	RADIUS, SSH, and Web
Internet edge	OSPF, BFD, IPsec, and BGP	RADIUS and SSH

Traditionally, many network administrators have tried to maintain a protect-router ACL / firewall filter template for each type of router, and this quickly becomes an operational burden because creating a customized firewall filter for each router is time consuming. The “protect-router” firewall filter becomes so complex, and so large, that it makes the management more difficult as the administrator scales the network, or adds new services. As networking services such as RA-DIUS servers or BGP peers change over time, the protect-router ACL / firewall filter needs to be updated to reflect these changes.

It's possible to create a simple security framework to support a wide variety of protocols and services by leveraging the power of the Junos operating system. This book illustrates simple building blocks based on commonly used protocols and services to make securing the routing engine an easy task: simply pick and choose what protocols and services are required for each router. Junos has configuration automation tools that enable dynamic prefix lists, and as your network scales and changes, the security framework automatically changes and adjusts without requiring additional input from the administrator.

Day One: Securing the Routing Engine on M/T/MX Series walks you through all the components and instructions on how to create a framework of firewall filters and policies to secure the routing engine, and in the final chapter you are presented with a final framework and configuration, which you can then implement on a router.

M/MX/T Routers

This book is geared towards the Juniper M/MX/T family of routers as they are most often deployed in very demanding environments. The amount of unwanted ingress traffic from the Internet is unbelievable. The router is constantly flooded with ICMP fragments, SSH login requests, file sharing protocols, and much more. In environments where all traffic must be treated as untrusted, it's essential that the routing engine is secured against unwanted traffic.

NOTE Juniper has newly released a high-performance, mid-range series of MX routers: MX5, MX10, and the MX40. While not used in the test bed for this book, everything in this book applies to the new devices. For more information about the new line of mid-range Juniper MX routers see <http://www.juniper.net/us/en/products-services/routing/mx-series/>.

Chapter 1

Firewall Filters

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The primary method of matching traffic and performing an action is performed with a firewall filter, and Chapter 1 details what firewall filters are and how to use them. A firewall filter can be applied to traffic in two directions: input and output. Input refers to traffic destined for the routing engine if applied to the control plane, or ingress traffic to a logical interface if applied to the data plane. Likewise, output refers to traffic sourced from the routing engine if applied to the control plane, or egress traffic from a logical interface if applied to the data plane.

This may seem a little confusing now, but once you walk through a few instances it becomes clear.

Firewall filters are used to match traffic and perform an action on the traffic matched, very similar to a *policy statement*. Firewall filters consist of one or more named terms and each term has a set of *match conditions*, and a set of *actions* or *non-terminating actions*. If no traffic is matched by any of the terms there is a hidden *implicit-rule* that discards all traffic. Here's the firewall filter configuration syntax:

```
[edit firewall family inet]
filter filter-name {
  term term-name {
    from {
      match-conditions;
    }
    then {
      action;
      nonterminating-actions;
    }
  }
  term implicit-rule {
    then discard;
  }
}
```

Firewall Families

Firewall filters are defined differently for different protocol family types and filters that are defined for a specific family type can only be applied to an interface using the same protocol family. For example, a firewall defined as `family inet` can only be applied to an interface under `family inet`. Table 1.1 lists some of the types of protocol families for firewall filters.

Table 1.1 Firewall Filter Protocol Family Types

Protocol Family	Description
any	Protocol-independent filter.
bridge	Protocol family bridge.
ccc	Protocol family circuit cross-connect.
inet	Protocol family IPv4.
inet6	Protocol family IPv6.
mpls	Protocol family MultiProtocol Label Switching.
vpls	Protocol family Virtual Private LAN Service.

NOTE This book focuses strictly on IPv4 or family `inet`. For more information about Junos and IPv6, see the Day One books about IPv6 at <http://www.juniper.net/dayone>.

How Firewall Filters are Evaluated

Firewall filters are evaluated from the top-down starting with the first term within the filter. If no matches are found within a term, the next term is evaluated until a match is found. If a match isn't found then the hidden *implicit-rule* at the end of the firewall filter discards the traffic.

Single-term Filters

If a firewall filter contains a single term, it is evaluated starting with the `from` statement. If there is a match, the action specified in the `then` statement is taken. If there is no match, the traffic is discarded. For example:

```
[firewall family inet]
filter example-filter {
  term single-term {
    from {
      source-address 10.0.0.0/8;
    }
    then {
      accept;
    }
  }
}
```

You can see in the firewall filter that there is a single term. If the packet has a source address of 10/8 it is accepted, and if there is no match, then the hidden implicit-rule discards the packet.

Multiple-term Filters

If the firewall filter has more than a single term, the filter is evaluated from the top-down starting with the first term. It's common to see term names in Junos documentation in ascending order such as `term 1`, `term 2`, and `term 3`. However, regardless of the name, the filter is always evaluated sequentially, top-down, until a match occurs. If there is a match, the action specified in the `then` statement is taken. If there is no match, the next term is evaluated. This process continues until there is a match or there are no more terms left to evaluate. If a packet passes all the way through a filter without matching any of the terms, the hidden *implicit-rule* at the end of the filter discards the packet.

If a packet matches the conditions in the `from` statement and the term doesn't contain a `then` statement, or doesn't have an *action* defined, the packet is accepted by default. Here's an example of a filter with several terms:

```
[firewall family inet]
filter example-filter {
  term 1 {
    from {
      source-address 10.0.0.0/8;
    }
    then {
      log;
      accept;
    }
  }
  term 2 {
    from {
      source-address 192.168.0.0/16;
    }
    then {
      accept;
    }
  }
  term 3 {
    from {
      source-address 172.16.0.0/12;
    }
  }
  term 4 {
    from {
      source-address 200.200.0.0/24;
    }
    then {
      log;
    }
  }
}
```

And in this firewall filter the packet is evaluated as illustrated in Figure 1.1.

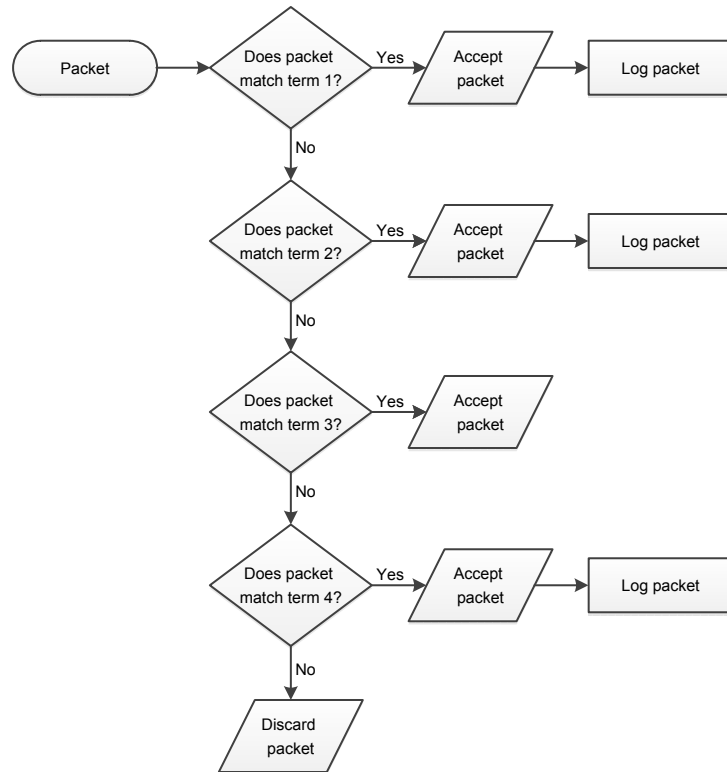


Figure 1.1 How A Packet is Evaluated in a Firewall Filter

BEST PRACTICE It's strongly recommended that you always configure an explicit action in the then statement. If there's no action, or the then statement is omitted completely, then packets that match the conditions in the from statement are permitted.

Firewall Filter Match Conditions

Match conditions are simply conditions for which the packet must match in order to be accepted, such as `protocol tcp` or `source-address 10.0.0.0/8`. Match conditions are specified in the from statement of the firewall filter. The order in which the match conditions appear isn't important, as all match conditions must be met in order

for the `then` statement to be taken. If there are no match conditions specified in a term, then all packets are matched.

When a single term has a list of the same match conditions, a match occurs if any one of the values in the list matches the packet, as shown here:

```
[firewall family inet]
filter example-filter {
  term 1 {
    from {
      protocol tcp;
      source-address 10.0.0.0/8;
      source-address 192.168.0.0/16;
      source-address 172.16.0.0/12;
    }
    then {
      accept;
    }
  }
}
```

Here the packet protocol is TCP, and if it has a source address of 10/8, 192.168/16, or 172.16/12 it will be accepted. It's important to remember that when a list of match conditions is defined, it requires only a single value to be matched within the list. In this example there are *two* match conditions that have to be met in order for the `from` statement to match: `protocol tcp` and the match list `source-address` which has three match conditions, but only one of those three match conditions has to match for the entire list to be evaluated as true.

ALERT Although only a single value must be matched within a list of match conditions, other match conditions outside of the list must be matched in addition to the list.

So in the firewall filter `example-filter`, the packet will be evaluated as depicted in Figure 1.2.

Firewall Filter Actions

Firewall actions fall into two groups: *terminating actions* and *non-terminating actions*, which are also called *action modifiers*. When a match occurs in the `from` statement, Junos can perform the specified action and/or non-terminating actions to the packet.

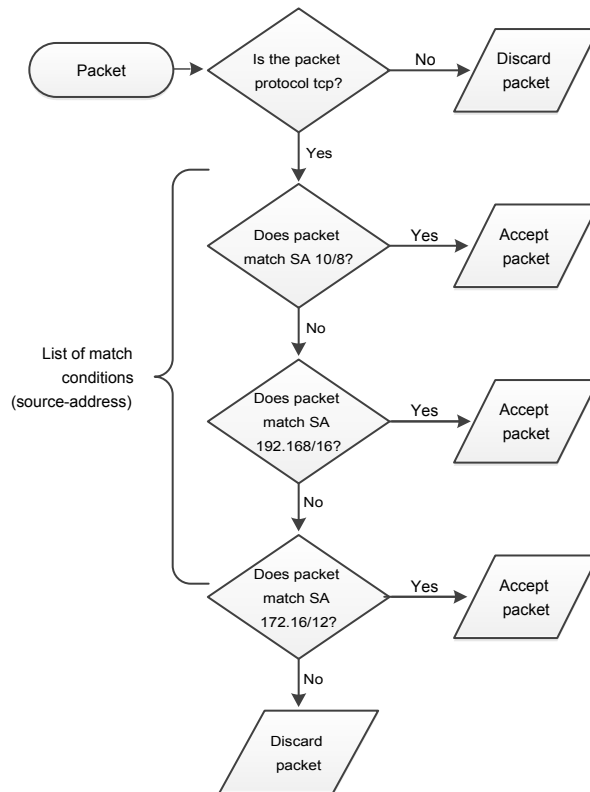


Figure 1.2 Flowchart of How the Firewall Filter Example is Evaluated

Here is the firewall filter configuration syntax for action and non-terminating actions:

```

filter filter-name {
  term term-name {
    from {
      match-conditions;
    }
    then {
      action;
      nonterminating-actions;
    }
  }
}

```

Actions

The action and non-terminating actions are defined within the `then` statement of a firewall filter. It's important to remember that if a match occurs and there isn't a defined `then` statement, the default action is to accept the packet. Only one action can be defined per term, but multiple non-terminating actions may be defined per term.

Table 1.2 Action Types

Action	Type	Description
<code>accept</code>	terminating	The packet is accepted.
<code>discard</code>	terminating	The packet is silently discarded.
<code>reject message-type</code>	terminating	The packet is rejected and the corresponding ICMP message is generated.
<code>routing-instance routing-instance</code>	terminating	Routes the packet using the specified routing-instance.
<code>next term</code>	non-terminating	Continues to the next term for evaluation.
<code>count counter-name</code>	non-terminating	Counts the number of packets passing this term. A counter name is specific to the filter that uses it, so all interfaces that use the same filter increment the same counter.
<code>forwarding-class class-name</code>	non-terminating	Classifies the packet to the specified forwarding class.
<code>log</code>	non-terminating	Logs the packet's header information in memory on the packet forwarding engine (PFE). This is a First In First Out (FIFO) buffer and is limited to about 400 entries.
<code>loss-priority priority</code>	non-terminating	Sets the scheduling priority of the packet.
<code>policer policer-name</code>	non-terminating	Applies rate limits to the traffic using the named policer.
<code>sample</code>	non-terminating	Samples the traffic on the interface. Use this non-terminating action only when sampling is enabled.
<code>syslog</code>	non-terminating	Records information in the system logging facility. This non-terminating action can be used with all actions except <code>discard</code> .

NOTE The action `next term` cannot be used with a terminating action in the same filter term, and can only be configured with non-terminating actions.

MORE? This book only deals with actions: `accept`, `reject`, `discard`, `count`, `log`, and `policer`. For more information on the other actions visit http://www.juniper.net/techpubs/en_US/junos10.4/topics/usage-guidelines/policy-configuring-actions-in-firewall-filter-terms.html.

The action `accept` is the most commonly used because the default action of a firewall filter is to `discard` all traffic that isn't matched. If the objective is to match certain packets and deny them, the most common action is `discard`. The `discard` action silently drops the packet without generating an ICMP message. This method has security benefits as the routing engine does not give any clues as to why the traffic was dropped. In certain situations it makes complete sense to match a packet and deny it and generate a corresponding ICMP message. The `reject` action discards the packet and generates and sends an ICMP `administratively-prohibited` message by default, but there are other ICMP messages than can be specified to override the default, such as:

- `administratively-prohibited` (this is the default)
- `bad-host-tos`
- `bad-network-tos`
- `fragmentation-needed`
- `host-prohibited`
- `host-unknown`
- `host-unreachable`
- `network-prohibited`
- `network-unknown`
- `network-unreachable`
- `port-unreachable`
- `precedence-cutoff`
- `precedence-violation`
- `source-host-isolated`
- `source-route-failed`
- `tcp-reset` (If the original packet was TCP, a TCP reset segment is generated, and if the original packet wasn't TCP, no response is generated.)

Here's an example firewall filter that matches packets with a RFC1918 source address and rejects the packets with an ICMP message of network-prohibited:

```
[firewall family inet]
filter example-filter {
  term 1 {
    from {
      source-address 10.0.0.0/8;
      source-address 192.168.0.0/16;
      source-address 172.16.0.0/12;
    }
    then {
      reject network-prohibited;
    }
  }
}
```

Non-terminating Actions

Non-terminating actions are additional processing that the router is able to perform on a packet. One or more non-terminating actions are able to be defined per term. The packets are still accepted, discarded, or rejected as defined by the corresponding action, but the non-terminating actions are also performed.

Counters

Counters are simply a tuple of the counter name, number of bytes passed, and total number of packets. If traffic matches a term and the non-terminating action count is called, the counter is incremented. Don't be afraid to use counters as they are processed in the packet forwarding engine (PFE) and operate at line rate without impacting performance. Here's an example:

```
[firewall family inet]
filter example-filter {
  term 1 {
    from {
      source-address 10.0.0.0/8;
      source-address 192.168.0.0/16;
      source-address 172.16.0.0/12;
    }
    then {
      count rfc1918;
      accept;
    }
  }
}
```

The example firewall filter illustrates how to define a counter. Use the non-terminating action `count` followed by a name for the counter.

NOTE Counter names are aggregated by default. If there are multiple firewall filters using the same counter name, each counter will increment an aggregated counter name. To change this behavior so that counter names are aggregated per interface, use the `interface-specific` knob within the firewall filter.

NOTE There can only be a single counter defined per term.

Counters give instant visibility into how the firewall filter is working and are often used to detect mistakes in a firewall filter or abnormal traffic conditions. A common use of counters is to count the number of bytes and packets that have been discarded at the end of a firewall filter.

Logging

If you use the non-terminating action `log`, it records the packet's header information and places it into a FIFO buffer on the PFE. This FIFO holds about 400 entries.

ALERT! The non-terminating action `log` is completely different from `syslog`. Although both `log` and `syslog` log the packet's header, the non-terminating action `log` operates at line-rate, storing the information in PFE memory, whereas the non-terminating action `syslog` sends the packet's header to the routing engine to be processed, which is subject to the internal routing engine policer.

Logging is a great tool to spot-check firewall filters without a loss in performance. It's commonly used to log traffic that wasn't matched in a firewall filter. Here's a firewall filter example that defines logging by using the non-terminating action `log`:

```
[firewall family inet]
filter example-filter {
  term 1 {
    from {
      source-address 10.0.0.0/8;
      source-address 192.168.0.0/16;
      source-address 172.16.0.0/12;
    }
    then {
      accept;
    }
  }
}
```

```

term 2 {
  then {
    log;
    discard;
  }
}

```

If the packet has a RFC1918 source address it will be accepted. All other packets are discarded and logged to the PFE.

ALERT! You should use logging sparingly as all logs are placed into an in-memory FIFO that only supports about 400 entries. If too much traffic is being logged, the FIFO rotates the logged traffic out so fast it isn't useful.

NOTE There can only be one log per firewall term.

Logging Information

The packet header is logged along with the exact time the packet was processed, and additional information such as the ingress interface and the firewall filter action is also logged. Table 1.3 depicts the types of logging information and their filter names.

Table 1.3 Types of Logging Information

Logging Information	Description
Time	Time that the event occurred.
Filter	Name of the filter that matched the traffic.
Act	Filter action:
	A – Accept
	D – Discard
	R – Reject
Interface	Ingress interface for the packet.
Protocol	Packet's protocol name.
Src address	Packet's source address and port.
Dest address	Packet's destination address and port.

Applying Firewall Filters

After a firewall filter has been created, it needs to be applied to an interface to become active. Only firewall filters that have been applied to

an interface are evaluated, and once applied, the filter is installed into the PFE and all processing is performed at line rate.

Direction

Firewall filters can be applied in two directions: *input* and *output*. Firewall filters applied with the direction of input match ingress traffic and firewall filters applied with the direction of output will match egress traffic.

To apply a firewall filter use the `filter` statement under the interface's protocol family, specify the direction, then the firewall filter name.

Input:

```
# set interfaces lo0.0 family inet filter input example-filter
```

Output:

```
# set interfaces lo0.0 family inet filter output example-filter
```

NOTE Be careful applying firewall filters to the interface `lo0.0`. The firewall filter must account for all management and routing protocols that are used on the router. It's recommended to always use the console or the command `commit confirmed` when applying firewall filters to the routing engine in case you accidentally lock yourself out.

How to Create a Firewall Filter

Let's put all the pieces together and try creating our own firewall filter:

1. Begin by going into `configure` mode on your router.

```
dhanks@MX80> configure  
Entering configuration mode
```

2. Now, create and name the firewall filter `example-filter` with a single-term of 1 with a non-terminating action of `log`:

```
[edit]  
dhanks@MX80# set firewall family inet filter example-filter term 1 then log
```

3. Now create a counter called `example-filter`.

```
[edit]  
dhanks@MX80# set firewall family inet filter example-filter term 1 then count example-filter
```

4. Create a terminating action of `accept`:

```
[edit]  
dhanks@MX80# set firewall family inet filter example-filter term 1 then accept
```

5. Double-check the firewall filter for accuracy:

```
[edit]
dhanks@MX80# show | compare
[edit]
+ firewall {
+   family inet {
+     filter example-filter {
+       term 1 {
+         then {
+           count example-filter;
+           log;
+           accept;
+         }
+       }
+     }
+   }
+ }
```

6. Apply the firewall filter `example-filter` to the loopback interface `lo0.0`:

```
[edit]
dhanks@MX80# set interfaces lo0.0 family inet filter input example-filter
```

7. Review all the changes for accuracy and commit:

```
dhanks@MX80# show | compare
[edit interfaces lo0 unit 0 family inet]
+   filter {
+     input example-filter;
+   }
[edit]
+ firewall {
+   family inet {
+     filter example-filter {
+       term 1 {
+         then {
+           count example-filter;
+           log;
+           accept;
+         }
+       }
+     }
+   }
+ }
```

```
[edit]
dhanks@MX80# commit and-quit
commit complete
Exiting configuration mode
```

```
dhanks@MX80>
```

ALERT! Always use the `commit confirmed 5` command when applying firewall filters to the control plane via `interface lo0.0`. This commits the configuration to the router and if the command `commit` isn't applied again within 5 minutes, the router automatically rolls back to the previous configuration. This way if you make a mistake you haven't locked yourself out.

Try It Yourself: Create Your Own Filter

Although the firewall filter just completed may not be particularly useful, it successfully demonstrates how to create a firewall filter with a single term and with both terminating and non-terminating actions. While Chapter 5 dives deep into firewall filters to protect the engine, including major routing protocols and management tools, try completing a few firewall filters now, for practice. Use the seven-step process and implement some of the terms, actions, and non-terminating actions discussed.

Firewall Filters: Data Plane versus Control Plane

Depending on where the firewall filter is applied, it will filter traffic on the data or control plane. If the firewall filter is applied to the `interface lo0.0` it filters all control plane traffic:

```
interfaces {
  lo0 {
    unit 0 {
      family inet {
        filter {
input accept-all;
        }
      address 127.0.0.1/32;
    }
  }
  ge-0/0/0
    unit 0 {
      family inet {
        address 10.0.0.1/30;
      }
    }
}
firewall {
  family inet {
    filter accept-all {
      term 1 {
        then {
          count accept-all;
          log;
          accept;
        }
      }
    }
  }
}
```

```

}
}
}
}
}

```

In this example, the firewall filter `accept-all` applied to interface `lo0.0` in the direction of input, counts, logs, and accepts all ingress traffic destined to the routing engine regardless of the source. For example, if an OSPF Hello Packet was received from interface `ge-0/0/0.0`, it would be accepted by the firewall filter `accept-all` because routing protocols such as OSPF require the routing engine to handle the Hello Packets.

A good way to think about firewall filters applied to interface `lo0.0` is that any traffic that is destined to, or has originated from, the routing engine is evaluated as depicted in Figure 1.3.

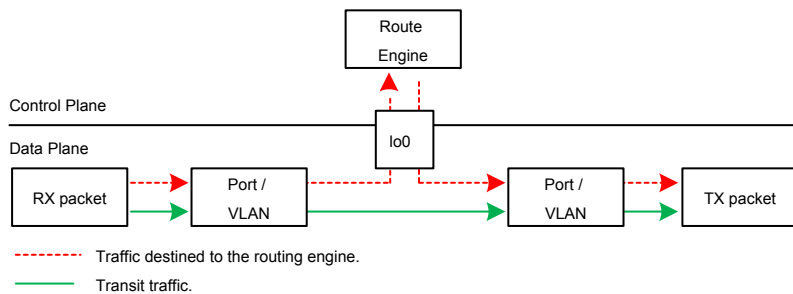


Figure 1.3 A Visual Representation of How Transit and Traffic Destined to the Routing Engine are Handled

Traffic destined to the router is always sent to the routing engine. These packets are evaluated by firewall filters applied on ingress logical interfaces first, and if the traffic is accepted it is then evaluated by firewall filters applied to the interface `lo0.0`. The reverse is true for packets sourced from the routing engine, such as a ping or telnet, to a remote host. This traffic is evaluated by the firewall filters applied on the interface `lo0.0`, and if the traffic is accepted it is then evaluated by firewall filters on the egress logical interface. When traffic is sourced from the routing engine, care needs to be taken to also allow return traffic.

Transit traffic that flows through the router doesn't need to be processed by the routing engine, therefore it isn't subject to firewall filters applied to the interface `lo0.0`. Transit packets are always evaluated by firewall

filters applied *on the ingress logical interface first*, and if accepted, will be subject to the applied egress filters.

Firewall Filter Chaining

The Junos operating system evaluates multiple firewall filters applied to an interface from left-to-right in the order they were configured. Each firewall filter is evaluated sequentially, top-down, until a match is found. When a match is found no more firewall filter terms or filters are evaluated. When Junos reaches the end of the firewall filter without a match, it moves to the next firewall filter in the list until there are no more firewall filters. This process is depicted in Figure 1.4.

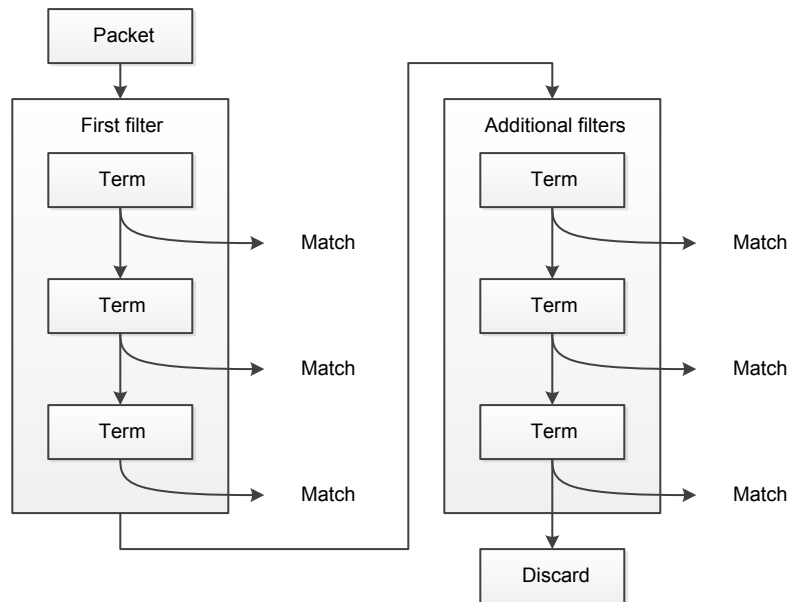


Figure 1.4 Visual Representation of How Firewall Chaining is Evaluated

Chaining firewall filters is a very powerful tool as it allows for a modular framework. Instead of creating a single firewall filter with numerous terms, it's possible to create individual firewall filters and chain them together.

```

interfaces {
  lo0 {
    unit 0 {
      family inet {
        filter {
  
```

```
        input-list [ accept-web accept-ssh discard-all ];
    }
    address 127.0.0.1/32;
}
}
}
}

firewall {
    family inet {
        filter accept-web {
            term 1 {
                from {
                    port web;
                }
                then {
                    accept;
                }
            }
        }
        filter accept-ssh {
            term 1 {
                from {
                    port ssh;
                }
                then {
                    accept;
                }
            }
        }
        filter discard-all {
            term 1 {
                then {
                    count discard-all;
                    log;
                    discard;
                }
            }
        }
    }
}
```

This example defines three firewall filters `accept-web`, `accept-ssh`, and `discard-all`. These firewall filters are applied to the interface `lo0.0` as an `input-list`, also known as a chain. The firewall filters will be evaluated left-to-right in the order they were entered in the `input-list`: `accept-web`, `accept-ssh`, and `discard-all`.

ALERT! Routing policy processing and firewall filter processing are completely different.

An important distinction when comparing firewall filter and routing policy processing is when the processing stops. In a routing policy chain the processing doesn't stop until a terminating action is met or a match isn't found. In a firewall filter chain the processing stops immediately after a match, even if you omit a terminating action.

ALERT! A maximum of 16 firewall filters can be applied as an `input-list` or `output-list`.

No matter how complex the problem, it's possible to break it down into simple building blocks that are able to be used elsewhere. Creating modular, individual firewall filters provides the added benefit of reusing them elsewhere in the configuration, thus allowing the operator to “do more with less” and not reinvent the wheel.

For a perfect example, take two completely different problems such as writing a firewall filter for BGP peering traffic to the Internet and writing a firewall filter to protect the routing engine. At a high level, the two tasks are not related, but once these items are broken down into simple building blocks, patterns and common tasks appear. Individually each filter would want to block ICMP fragments and eventually discard that traffic at the end of the policy, but by creating common firewall filters to specifically perform these tasks and chaining them together, they can work together.

TIP Common firewall filters should be focused enough to accomplish the immediate task at hand, but generic enough to work in any scenario. Then you have a collection of common firewall filters as a bag of building blocks, and each block can be built on top of others in different ways for different filtering goals. It's like network masonry.

Nested Firewall Filters

It's possible to create a filter that references another filter. The only limitation is that this can only be done once per term and doesn't support recursive firewall filters. Consider the following example:

```
[edit firewall family inet]
filter accept-web {
  term 1 {
    from {
      port web;
    }
    then {
      accept;
    }
  }
}
```

```
    }  
  }  
  filter accept-ssh {  
    term 1 {  
      from {  
        port ssh;  
      }  
      then {  
        accept;  
      }  
    }  
  }  
  
  filter accept-web-ssh {  
    term accept-web {  
      filter accept-web;  
    }  
    term accept-ssh {  
      filter accept-ssh;  
    }  
  }  
}
```

In this sample firewall filter two regular filters are defined: `accept-web` and `accept-ssh`. Each of these filters looks for packets using the specified port and accepts the traffic. These two firewall filters can be combined into a single filter called `accept-web-ssh`, which allows traffic from both filters `accept-web` and `accept-ssh`.

Because there is a hard limit to the number of firewall filters that can be applied in a chain, using nested firewall filters is a great tool to combine common filters and have them counted as a single filter.

Summary

Firewall filters are the building blocks used to match and take action on traffic. It's critical to understand how firewall filters are evaluated in order to build a strong line of defense. There are many forms of firewall filters: single-term, multiple-term and nested filter. This flexibility gives you the ability to solve complex problems with ease.

Writing firewall filters is a fine balance of science and art. Technically the filter needs to match traffic and take action. Artistically the filters need to be written so that each filter can be reused for other tasks. When the filters are chained together, careful thought needs to be taken to ensure that the desired result is achieved.

Chapter 2

Policers

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Chapter 2 explains what policers are and shows you how to use them. Rate limiting is performed via a firewall action that references a policer, and firewall filters and policers are used together to match traffic and enforce a rate limit. Let's start.

Policers Overview

Policing is synonymous with *rate limiting*. Policers work together with firewall filters to set bandwidth restrictions on matched traffic and enforce consequences on traffic exceeding the bandwidth limitation defined by the policer.

Despite the constant confusion, policers are very simple to configure. There are three major components that need to be defined: policer name, *if-exceeding* parameters, and the action. The policer configuration syntax is as follows:

```
firewall {
  policer policer-name {
    filter-specific;
    if-exceeding {
      bandwidth-limit bps;
      bandwidth-percent number;
      burst-size-limit bytes;
    }
    then {
      policer-action;
    }
  }
}
```

Token Bucket Algorithm

Junos policers use the *token-bucket algorithm* to enforce an average bandwidth limit over time, while allowing for bursts that exceed the bandwidth limit.

Packets arrive at the policer at the rate at which they were originally sent and the policer decision process is run for each ingress packet. There are three parameters that the policer uses to determine if the packet will be accepted: *packet size*, *bandwidth-limit*, and *burst-size-limit*. The policer simply subtracts the packet size from the current size

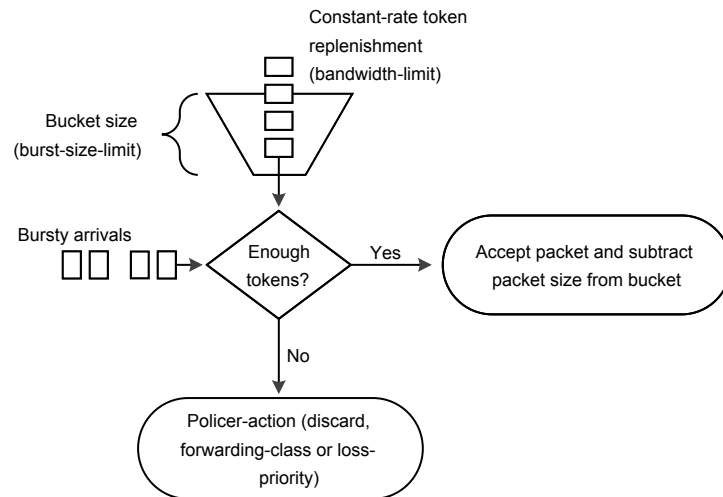


Figure 2.1 Visual Representation of the Token-bucket Algorithm

of the token bucket – if the result is greater than zero the packet is accepted. If there aren't enough tokens in the bucket, then the policer handles the packet in one of two ways: *hard* or *soft policing*.

- Hard policing: the packet will simply be discarded.
- Soft policing: the packet can be marked two ways.
 - Set the packet loss priority (PLP).
 - Set the forwarding class.

As packets are accepted they leave the policer at the rate at which they were originally sent, which is why the token-bucket algorithm is better suited for bursty traffic.

NOTE A good way to think about the token-bucket algorithm versus the leaky-bucket algorithm is that the token bucket results in bursty departures and the leaky bucket results in smooth departures.

MORE? For more information on the token-bucket algorithm see the book, *QOS-Enabled Networks* (by Barreiros & Lundqvist, Wiley & Sons, 2011) at www.juniper.net/books.

Bandwidth-limit

The `bandwidth-limit` Junos parameter is the knob that adjusts how many bits (on average) are allowed during a one second interval. This is referred to as *bits per second* (bps), but also accepts more human-readable values such as k (1000), m (1,000,000), and g (1,000,000,000). Valid ranges are 32 Kbps through 40 Gbps.

For example, to create a `bandwidth-limit` of 100 megabits per second (Mbps) use the value `100m`, like this:

```

policer example-policer {
  if-exceeding {
    bandwidth-limit 100m;
    burst-size-limit 625k;
  }
  then {
    discard;
  }
}

```

TIP The `bandwidth-limit` knob is calculated in bits per second (bps).

The nature of the token-bucket algorithm allows for small bursts of traffic above the defined `bandwidth-limit`, but still enforces an average bandwidth-limit over the time period of one second. The two graphs in Figure 2.2 illustrate how the bandwidth-limit is enforced, but still allow for bursts beyond the actual bandwidth-limit. These illustrations assume the physical interface speed is 100m and the bandwidth-limit has been set to 50m.

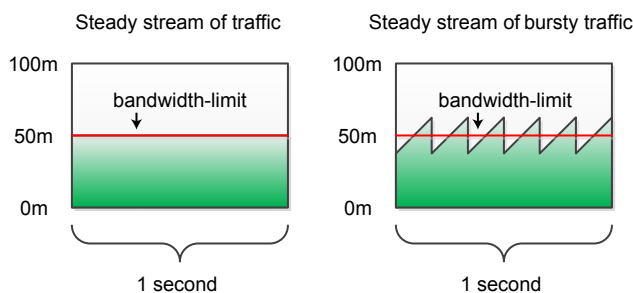


Figure2.2 Example of Bandwidth-limit and Traffic Bursting

NOTE These graphs do not illustrate real-world traffic and are not drawn to scale, but hopefully they act as a visual aid to help you understand how traffic is able to burst beyond the bandwidth-limit, even though the

average bandwidth over a period of 1 second is still equal to the bandwidth-limit.

In the first graph there is a steady stream of traffic using exactly 50Mbps. In this case the token bucket remains full because as traffic comes in, the bandwidth-limit replenishes tokens into the bucket at the exact same rate as the incoming traffic.

The second graph illustrates a steady stream of traffic around 40 Mbps; it has small bursts up to 60 Mbps, but then drops back down to 40 Mbps. The token bucket is nearly depleted after each burst, but the traffic drops back down to 40 Mbps to allow the tokens to fill up the bucket so there are enough tokens to burst again.

TIP The `bandwidth-limit` defines how fast the token bucket will be refilled every clock cycle. If the current traffic is below the `bandwidth-limit`, this means a surplus is created in the token bucket, whereas if the current traffic is above the `bandwidth-limit`, there is a deficit in the token bucket, until it's depleted and begins policing traffic.

Burst-size-limit

The `burst-size-limit` defines how large the token bucket is. The token bucket controls the amount of burst permitted before packets are policed. Since the `bandwidth-limit` is a constant value used to enforce an average bandwidth over time, the only variable to adjust how much burst the policer will allow is `burst-size-limit`.

There are two common methods for calculating the `burst-size-limit`.

1. If the average packet size is known in advance, set the `burst-size-limit` to average packet size * 10. If the packet size is unknown, the only option is to set the `burst-size-limit` to the physical interface's MTU * 10.
2. Set the `burst-size-limit` to the amount of data the physical interface can transmit within a 5ms window.

NOTE The author finds that setting the `burst-size-limit` to the amount of bandwidth the physical interface can transmit in 5ms is the best method.

Setting a smaller `burst-size-limit` for real-time traffic is recommended, but setting the value too low could cause the token bucket to

constantly be empty. Setting a larger `burst-size-limit` allows for larger bursts, but traffic would not be policed. The most common and preferred method is using the physical interface calculation, such as these common physical interfaces and burst sizes:

Interface	burst-size-limit
100Mbps (fe)	62,500 bytes.
1Gbps (ge)	625,000 bytes.
10Gbps (xe)	6,250,000 bytes.

ALERT! Be careful. `burst-size-limit` is calculated in bytes, while `bandwidth-limit` is calculated in bits per second.

It's easy to forget to convert the bits to bytes and miscalculate the `burst-size-limit`. Here's a best practice `burst-size-limit` formula:

$$(\text{interface bandwidth [bps]} / 1000 [\text{ms}]) * 5\text{ms} / 8 = \text{burst-size-limit [bytes]}$$

And an example of the formula for 1Gbps interface:

$$1,000,000,000 \text{ bps} / 1000 \text{ ms} = 1,000,000$$

$$1,000,000 * 5\text{ms} = 5,000,000$$

$$5,000,000 \text{ bps} / 8 = 625,000 \text{ bytes}$$

$$\text{burst-size-limit} = 625,000 \text{ bytes}$$

NOTE The link between the PFEs and the routing engine is 1Gbps. All policers defined in this book will have a `burst-size-limit` of 625,000 bytes.

Rate-limiting Traffic

Firewall filters and policers work together to match traffic and apply rate limiting, enabling you to have exact control of which traffic should be rate-limited and which traffic is allowed to pass through unrestricted. Rate-limiting traffic requires two steps:

- Create a policer with rate limiting parameters.
- Create a firewall filter to match the traffic that needs to be policed and use the `policer` action-modifier.

How to Rate Limit with Both a Policer and a Firewall Filter

To rate limit only SSH traffic and accept all other traffic to the routing engine there needs to both a policer and a firewall filter.

1. Let's start by creating a policer to rate limit traffic to 10 Mbps:

```
[edit]
dhanks@MX80# set firewall policer 10m if-exceeding bandwidth-
limit 10m burst-size-limit 625000
```

```
[edit]
dhanks@MX80# set firewall policer 10m then discard
```

2. Now create a firewall filter to match SSH traffic:

```
[edit]
dhanks@MX80# set firewall family inet filter limit-ssh term 1
from protocol tcp port ssh
```

3. Reference the policer 10m as the action-modifier and accept:

```
[edit]
dhanks@MX80# set firewall family inet filter limit-ssh term 1
then policer 10m accept
```

4. Don't forget to add a final term to accept all other control plane traffic, but not do police it:

```
[edit]
dhanks@MX80# set firewall family inet filter limit-ssh term 2
then accept
```

5. Apply the limit-ssh firewall filter to the interface lo0 with the direction of input:

```
[edit]
dhanks@MX80# set interfaces lo0.0 family inet filter input
limit-ssh
```

6. The final step is to double-check the configuration and apply the new configuration:

```
dhanks@MX80# show | compare
[edit interfaces lo0 unit 0 family inet]
+   filter {
+     input limit-ssh;
+   }
[edit]
+ firewall {
+   family inet {
+     filter limit-ssh {
+       term 1 {
+         from {
+           protocol tcp;
+           port ssh;
+         }

```

```

+         then {
+             policer 10m;
+             accept;
+         }
+     }
+     term 2 {
+         then accept;
+     }
+ }
+ }
+ policer 10m {
+     if-exceeding {
+         bandwidth-limit 10m;
+         burst-size-limit 625k;
+     }
+     then discard;
+ }
+ }

```

```

[edit]
dhanks@MX80# commit and-quit
commit complete
Exiting configuration mode

```

```
dhanks@MX80>
```

The first term in the filter specifically matches SSH traffic by looking at the port and protocol. If a match is found, it will police SSH traffic using the `policer 10m`. The second term is there to override the implicit discard all at the end of the policer. The simple firewall filter and policer illustrate how it's possible to only police SSH traffic while allowing all other traffic to the routing engine.

Filter-specific Versus Term-specific

By default, the Junos operating system creates a policer instance per firewall term. For example, if a firewall contained three terms and two terms had an action-modifier of `policer 10m`, there would be two instances of the policer. This allows each term in a firewall filter to have its own policer so that it doesn't have to share a policer with other terms, such as shown here:

```

firewall {
    family inet {
        filter limit-ssh-http {
            term 1 {
                from {
                    protocol tcp;
                    port ssh;
                }
            }
        }
    }
}

```

```
    }
    then {
        policer 10m;
        accept;
    }
}
term 2 {
    from {
        protocol tcp;
        port http;
    }
    then {
        policer 10m;
        accept;
    }
}
term 3 {
    then accept;
}
}
}
policer 10m {
    if-exceeding {
        bandwidth-limit 10m;
        burst-size-limit 625k;
    }
    then discard;
}
}
```

In this example, the firewall filter `limit-ssh-http` has a total of three terms. Terms 1 and 2 reference the same policer. Although both terms reference the same policer, each term receives its own instance of the policer, guaranteeing that SSH and HTTP traffic will be individually policed at 10 Mbps. Use the `show firewall filter` command to verify the number of policer instances, such as shown here:

```
dhanks@MX80> show firewall filter limit-ssh-http
```

```
Filter: limit-ssh-http
Policers:
Name                               Packets
10m-1                               0
10m-2                               0
```

```
dhanks@MX80>
```

The firewall filter `limit-ssh-http` has two policer instances: `10m-1` and `10m-2`. The first instance is assigned to the first term referencing the policer and the second instance is assigned to the second term referenc-

ing the policer.

There may be scenarios when it's desirable to create a policer that can be shared across all the terms within a firewall filter. Use the `filter-specific` knob inside of the policer to change how Junos calculates the number of instances to create per filter. When a policer is defined using the `filter-specific` knob, only one instance of the policer is created per firewall filter and traffic is aggregated across the terms.

```
firewall {
  family inet {
    filter limit-ssh-http {
      term 1 {
        from {
          protocol tcp;
          port ssh;
        }
        then {
          policer 10m;
          accept;
        }
      }
      term 2 {
        from {
          protocol tcp;
          port http;
        }
        then {
          policer 10m;
          accept;
        }
      }
      term 3 {
        then accept;
      }
    }
  }
  policer 10m {
    filter-specific;
    if-exceeding {
      bandwidth-limit 10m;
      burst-size-limit 625k;
    }
    then discard;
  }
}
```

Here, the firewall filter `limit-ssh-http` has three terms, two of which have an action-modifier of `policer 10m`. The difference is that the `policer 10m` now has the knob `filter-specific`, so there is only one policer instance created for the firewall filter `limit-ssh-http`:

```
dhanks@MX80> show firewall filter limit-ssh-http
```

```
Filter: limit-ssh-http
```

```
Policers:
```

Name	Packets
10m	0

```
dhanks@MX80>
```

The firewall filter `limit-ssh-http` now only has a single policer instance. Both SSH and HTTP traffic will be policed to 10 Mbps as a whole.

Summary

Historically, policing has been one of the more difficult topics to understand in depth. Stepping back from all of the implementation details and focusing on the actual token bucket algorithm helps you to understand how policing actually works. The bandwidth limit and token bucket size work together to enforce how much traffic is allowed during a one second interval.

If traffic exceeds the limits defined by the policer, the traffic is either soft policed or hard policed. Soft policing traffic allows you to set the PLP or forwarding class of the packet. Hard policing will silently discard the traffic.

Policers come in two flavors: term-specific and filter-specific. By default policers are term-specific. A term-specific policer will have an instance created for every term in a firewall filter for which it's applied. On the other hand, if a policer is filter-specific, regardless of the number of terms to which the policer is applied, only a single instance of the policer is created.

Chapter 3

Viewing Counters, Logs, and Policers

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When firewall filters are applied to the routing engine, it's important for the administrator to know if the filters are working as expected or if there is traffic being discarded unintentionally. There are three tools covered in this chapter that can be used to receive feedback on firewall filters: counters, logs, and policer counters.

Fire up your router and follow along. Let's get at it.

Viewing Firewall Filter Counters

Counters are named differently depending on how you apply the firewall filter and if the firewall filter has the `interface-specific` knob enabled. When firewall filters are applied to an interface using `filter input` or `filter output` the counter names will be aggregated.

The following firewall filter accepts TCP and UDP traffic and counts it.

```
filter accept-tcp-udp {
  apply-flags omit;
  term accept-all-tcp {
    from {
      protocol tcp;
    }
    then {
      count accept-all-tcp;
      accept;
    }
  }
  term accept-all-udp {
    from {
      protocol udp;
    }
    then {
      count accept-all-udp;
      accept;
    }
  }
}
```

The firewall filter `accept-tcp-udp` will be applied to two separate interfaces `ge-1/3/9.0` and `lo0.0`.

```
ge-1/3/9 {
  unit 0 {
    family inet {
      filter {
        input accept-tcp-udp;
      }
    }
  }
}
```

```

    }
    address 1.1.1.1/30;
  }
}
lo0 {
  unit 0 {
    family inet {
      filter {
        input accept-tcp-udp;
      }
      address 172.16.2.1/32;
    }
  }
}

```

Notice although the firewall filter is applied to two separate interfaces, there is only a single instance of the counter.

```
dhanks@MX80> show firewall filter accept-tcp-udp
```

```
Filter: accept-tcp-udp
```

```
Counters:
```

Name	Bytes	Packets
accept-all-tcp	222752	3777
accept-all-udp	1088	14

```
dhanks@MX80>
```

There are two ways to change this behavior. Enabling the interface-specific knob on the firewall filter `accept-tcp-udp` or applying the filter using a `filter input-list`.

interface-specific

Let's take a look at the interface-specific knob and how it changes the behavior.

```
dhanks@MX80> configure
```

```
Entering configuration mode
```

```
dhanks@MX80# set firewall family inet filter accept-tcp-udp interface-specific
```

```
[edit]
```

```
dhanks@MX80# commit and-quit
```

```
commit complete
```

```
Exiting configuration mode
```

```
dhanks@MX80> show firewall filter accept-tcp-udp-lo0.0-i
```

```
Filter: accept-tcp-udp-lo0.0-i
```

```
Counters:
```

Name	Bytes	Packets
accept-all-tcp-lo0.0-i	12736	243
accept-all-udp-lo0.0-i	0	0

```
dhanks@MX80> show firewall filter accept-tcp-udp-ge-1/3/9.0-i
```

```
Filter: accept-tcp-udp-ge-1/3/9.0-i
```

```
Counters:
```

Name	Bytes	Packets
accept-all-tcp-ge-1/3/9.0-i	0	0
accept-all-udp-ge-1/3/9.0-i	0	0

```
dhanks@MX80>
```

Notice how the counters are named differently. The counter naming convention when using the interface-specific knob is <counter_name>-<interface>.<unit>-<direction>. Also note that the firewall filter name has changed as well.

Firewall Chaining

Now let's try applying the firewall filter as a chain.

```
dhanks@MX80> configure
```

```
Entering configuration mode
```

```
[edit]
```

```
dhanks@MX80# delete interfaces lo0.0 family inet filter
```

```
[edit]
```

```
dhanks@MX80# set interfaces lo0.0 family inet filter input-list accept-tcp-udp
```

```
[edit]
```

```
dhanks@MX80# delete interfaces ge-1/3/9.0 family inet filter
```

```
[edit]
```

```
dhanks@MX80# set interfaces ge-1/3/9.0 family inet filter input-list accept-tcp-udp
```

```
[edit]
```

```
dhanks@MX80# commit and-quit
```

```
commit complete
```

```
Exiting configuration mode
```

```
dhanks@MX80> show firewall filter lo0.0-i
```

```
Filter: lo0.0-i
```

```
Counters:
```

Name	Bytes	Packets
accept-all-tcp-lo0.0-i	8600	150
accept-all-udp-lo0.0-i	0	0

```
dhanks@MX80> show firewall filter ge-1/3/9.0-i
```

```
Filter: ge-1/3/9.0-i
```

```
Counters:
```

Name	Bytes	Packets
accept-all-tcp-ge-1/3/9.0-i	0	0
accept-all-udp-ge-1/3/9.0-i	0	0

```
dhanks@MX80>
```

Notice that when using firewall chaining, the counter names are identical, but the firewall filter name is different than using the `interface-specific` method. The filter naming convention for firewall chaining is `<interface>.<unit>-<direction>`.

Table 3.1 Filter Naming Convention

Method	Counter Naming Convention	Filter Naming Convention
Default	Counter Name	Filter Name
interface-specific	<counter_name>- <interface>.<unit>- <direction>	<filter_name>- <interface>.<unit>-<direction>
Firewall Chaining	<counter_name>- <interface>.<unit>- <direction>	<interface>.<unit>-<direction>

It's recommended that firewall filters be applied as a chain. The counter and firewall filter names are easy to read and remember. The data is also broken out per interface so that there is more information available during troubleshooting and analysis.

Naming Convention

It's recommended that you have a naming convention when creating firewall filters, terms, and counters. For purposes of this book, the convention is to keep the term and counter name the same when creating filters.

Let's take a look at a firewall filter that has multiple terms and you can see how it varies from a firewall filter with a single term. In the next filter example, the firewall filter `accept-icmp` has two terms: `no-icmp-fragments` and `accept-icmp`. According to the convention of this book, each term has a counter with the same name as the term:

```
firewall {
  family inet {
```

```

filter accept-icmp {
  term no-icmp-fragments {
    from {
      is-fragment;
      protocol icmp;
    }
    then {
      count no-icmp-fragments;
      log;
      discard;
    }
  }
  term accept-icmp {
    from {
      protocol icmp;
      icmp-type [ echo-reply echo-request time-exceeded
unreachable source-quench router-advertisement parameter-
problem ];
    }
    then {
      policer management-5m;
      accept;
    }
  }
}
}

```

You can see that multiple term firewall filters follow the same logic as single term firewall filters, and you can use the `show firewall counter` command using the filter name and counter name as arguments:

```
dhanks@MX80> show firewall counter no-icmp-fragments-100.0-i filter 100.0-i
```

```
Filter: 100.0-i
```

```
Counters:
```

Name	Bytes	Packets
no-icmp-fragments-100.0-i	13824	12

```
dhanks@MX80> show firewall counter accept-icmp-100.0-i filter 100.0-i
```

```
Filter: 100.0-i
```

```
Counters:
```

Name	Bytes	Packets
accept-icmp-100.0-i	20312	213

```
dhanks@MX80>
```

Keeping the counter names in sync with the firewall filter's term names makes it easy to understand where the counter was applied.

Viewing the Firewall Filter Log

The second method for receiving feedback from your firewall filters is viewing the firewall filter log. This is a bit more straightforward than counters as the logs are located in the PFE memory for all filters, and since the firewall filters logs are stored in a central place, there's no distinction between them and there's no need to specify filter names or any other additional command-line arguments. Let's view an example firewall filter log by using the `show firewall log` command:

```
dhanks@MX80> show firewall log
```

```
Log :
```

```
Time      Filter  Action Interface Protocol  Src Addr      Dest Addr
03:38:36  lo0.0-i D      fxp0.0  ICMP     172.16.1.100 172.16.1.11
03:38:36  lo0.0-i D      fxp0.0  ICMP     172.16.1.100 172.16.1.11
03:38:36  lo0.0-i D      fxp0.0  ICMP     172.16.1.100 172.16.1.11
03:38:36  lo0.0-i D      fxp0.0  ICMP     172.16.1.100 172.16.1.11
03:38:31  lo0.0-i D      fxp0.0  ICMP     172.16.1.100 172.16.1.11
03:38:31  lo0.0-i D      fxp0.0  ICMP     172.16.1.100 172.16.1.11
```

You can see that the packet header is logged along with the exact time the packet was processed. Additional information such as the ingress interface and the firewall filter action is also logged.

Table 3.2 Firewall Filter Log Detail

Logging Information	Description
Time	Time that the event occurred.
Filter	Name of the filter that matched the traffic.
Act	Filter action:
	A – Accept
	D – Discard
	R – Reject
Interface	Ingress interface for the packet.
Protocol	Packet's protocol name.
Src address	Packet's source address and port.
Dest address	Packet's destination address and port.

Remember that the firewall log can only contain about 400 entries at any given time and is kept as a FIFO memory database located in the PFE. This allows log entries to be written at line-rate, but the drawback is that it's limited in scope.

The best practice calls for only using logging in scenarios where traffic shouldn't exist, such as a final discard firewall filter, as shown here:

```
firewall {
  family inet {
    filter discard-all {
      term discard-tcp {
        from {
          protocol tcp;
        }
        then {
          count discard-tcp;
          log;
          discard;
        }
      }
    }
    term discard-netbios {
      from {
        protocol udp;
        destination-port 137;
      }
      then {
        count discard-netbios;
        discard;
      }
    }
    term discard-udp {
      from {
        protocol udp;
      }
      then {
        count discard-udp;
        log;
        discard;
      }
    }
    term discard-icmp {
      from {
        protocol icmp;
      }
      then {
        count discard-icmp;
        log;
        discard;
      }
    }
    term discard-unknown {
      then {
        count discard-unknown;
        log;
        discard;
      }
    }
  }
}
```



```

    }
  }
}

```

When traffic reaches the last firewall filter `discard-all`, detailed information is available about that traffic in the firewall log.

To view additional logging detail, use the `show firewall log detail` command:

```

dhanks@MX80> show firewall log detail
Time of Log: 2011-03-27 03:38:36 UTC, Filter: lo0.0-i, Filter action: discard, Name of
interface: fxp0.0
Name of protocol: ICMP, Packet Length: 0, Source address: 172.16.1.100, Destination
address: 172.16.1.11
ICMP type: 0, ICMP code: 0
Time of Log: 2011-03-27 03:38:36 UTC, Filter: lo0.0-i, Filter action: discard, Name of
interface: fxp0.0
Name of protocol: ICMP, Packet Length: 0, Source address: 172.16.1.100, Destination
address: 172.16.1.11
ICMP type: 0, ICMP code: 0

```

As you can see, additional attributes such as packet length, source and destination ports, and IGMP types and codes are all available in the detailed view.

Viewing Firewall Policers

The third method of feedback from firewall filters are policers. When a policer is referenced in a firewall filter, a counter is automatically created using the policer name, interface name, unit number, and direction as you can see here:

```

interfaces {
  lo0 {
    unit 0 {
      family inet {
        filter {
          input-list [ accept-icmp accept-ssh discard-all ];
        }
        address 127.0.0.1/32;
      }
    }
  }
}
firewall {
  family inet {

```

```

filter accept-ssh {
  term accept-ssh {
    from {
      source-prefix-list {
        rfc1918;
      }
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      protocol tcp;
      destination-port ssh;
    }
    then {
      policer management-5m;
      count accept-ssh;
      accept;
    }
  }
}
}
}

```

To view the policer statistics, use the same `show firewall counter` command referencing the policer name, interface name, unit number, and direction:

```

dhanks@MX80> show firewall counter management-5m-1o0.0-i filter 1o0.0-i
Filter: 1o0.0-i
Policers:
Name                               Packets
management-5m-1o0.0-i             246

dhanks@MX80>

```

Summary

Voilà. There you have it. Three excellent ways to view what your firewall filters are doing. Counters count the number of bytes and packets that were matched. Policers count the number of bytes and packets that were policed. And the firewall log provides detailed information on packets that were matched and logged.

Chapter 4

Junos Configuration Automation

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Firewall filters and prefix lists are so tightly coupled that a single change can cause ripple effects throughout the entire configuration. As your configuration grows with new firewall filters and prefix lists, it poses ongoing operational challenges. Chapter 4 introduces you to Junos configuration automation, and is well worth the time it takes to review it.

BGP neighbor changes are a good example of complexity. If prefix lists are being used, the administrator must ensure that the BGP neighbor information is replicated into the prefix lists. If prefix lists aren't being used, the administrator has to go through each firewall filter and make sure that the BGP neighbor information is included in the firewall filter's `from` statement as shown below:

```
filter accept-bgp {
  term accept-bgp {
    from {
      source-address {
        192.0.2.5;
        192.0.2.6;
        192.0.2.7;
        192.0.2.8;
        192.0.2.9;
        192.0.2.10;
        192.0.2.11;
      }
      destination-address {
        10.255.255.5;
        10.255.255.6;
        10.255.255.7;
        10.255.255.8;
        10.255.255.9;
      }
      protocol tcp;
      port bgp;
    }
    then {
      count accept-bgp;
      accept;
    }
  }
}
```

Several different tools are available to mitigate the complexity and number of touch points of your Junos configuration. For all the tools covered in this chapter, the general rule of thumb is to let Junos do all the repetitive work. You won't find it in any data sheet but it's one of the reasons you choose Junos in the first place, right?

Apply-path

As your Junos configuration grows larger you'll see a lot of redundant information such as IP addresses and prefix lists. To help fight against this redundant information clouding up your configuration, there's a Junos feature called *apply-path*. Using keywords to match patterns inside of the Junos configuration, the *apply-path* feature applies the matched results into a dynamic prefix list.

The example below shows an IPv4 interface configuration and the corresponding static prefix list:

```
interfaces {
  ge-1/0/8 {
    unit 0 {
      family inet {
        address 10.0.8.6/30;
      }
    }
  }
  ge-1/0/9 {
    unit 0 {
      family inet {
        address 10.0.8.9/30;
      }
    }
  }
  ge-1/1/6 {
    unit 0 {
      family inet {
        address 10.0.2.1/30;
      }
    }
  }
  ge-1/1/7 {
    unit 0 {
      family inet {
        address 10.0.2.9/30;
      }
    }
  }
}
policy-options {
  prefix-list router-manual-ipv4 {
    10.0.2.0/30;
    10.0.2.8/30;
    10.0.8.4/30;
    10.0.8.8/30;
  }
}
```

You can see that the IPv4 addresses are originally defined on the interfaces, but to build a prefix list that contains all of the IPv4 addresses, there is a lot of keying in of redundant information. Creating a prefix list and manually adding one address at a time is mundane. As your network grows and changes, these prefix lists need to be updated as well. Static configuration is a risk as it could lead to the prefix list update being overlooked during an interface change, which would cause downtime or additional work to debug and resolve the issue.

Using `apply-path` allows Junos to do the repetitive work and maintain an up-to-date prefix list. The feature works by matching keywords inside the configuration using wildcards, then applying the matching addresses to create a dynamic prefix list. Let's try one:

```
prefix-list router-ipv4 {
  apply-path "interfaces <*> unit <*> family inet address <*>";
}
```

You can see in the `prefix-list router-ipv4` example that Junos begins at the `interfaces` stanza and using the `<*>` wildcard matches all interfaces, then matches all unit numbers, then explicitly moves into `family inet address` and matches all IPv4 addresses again with the `<*>` wildcard.

To see how Junos applied the wildcards to build the dynamic prefix list, use the `display inheritance` command when viewing the configuration:

```
[edit]
dhanks@MX80# show policy-options prefix-list router-ipv4 | display inheritance
##
## apply-path was expanded to:
## 10.0.8.4/30;
## 10.0.8.8/30;
## 10.0.2.0/30;
## 10.0.2.8/30;
##
apply-path "interfaces <*> unit <*> family inet address <*>";
```

```
[edit]
dhanks@MX80#
```

Now you can see that the prefix list `router-ipv4` contains four entries that match all of the IPv4 addresses on the physical interfaces.

ALERT! It's critical to understand that `apply-path` automatically fills out the prefix list using the *Direct* route instead of the *Local* route.

A Direct route in Junos is simply the logical interface's network address

plus the network mask. For example, if the logical interface's IP address was family inet address 10.0.0.1/30 the network address would be 10.0.0.0 and the network mask would be 30 bits, resulting in a Direct route of 10.0.0.0/30.

A Local route in Junos is the logical interface's IP address combined with a 32 bit mask. A good way to think of a Local route is as the actual IP address that resides on the router. For example if the logical interface IP address was family inet address 10.0.0.1/30 the Local route would be 10.0.0.1/32.

NOTE Using wildcards with apply-path is similar to the way that wildcards work with Junos configuration groups. For more information about Junos wildcards and configuration groups visit http://www.juniper.net/techpubs/en_US/junos10.4/topics/concept/junos-cli-wildcard-characters-configuration-groups-usage.html

The last match in your apply-path statement must be an IP address. For example if the last match was not an IP address, Junos would instantly return with an error on the CLI:

```
dhanks@MX80# set policy-options prefix-list test apply-path "interfaces <*>"
error: 'interface_name' is not IP address type
error: invalid value: interfaces <*>
```

```
[edit]
dhanks@MX80#
```

The apply-path feature can also be used to match IP addresses within routing protocols and services. Let's use an example matching all of the BGP neighbors and creating a prefix list called *bgp-neighbors*:

```
[edit]
dhanks@MX80# show protocols bgp
group AS65000 {
  type external;
  local-as 65000;
  neighbor 10.0.3.3;
  neighbor 10.0.3.4;
}
group AS65001 {
  type internal;
  local-address 10.0.6.1;
  neighbor 10.0.9.6;
  neighbor 10.0.9.7;
}
```

```
[edit]
dhanks@MX80#
```

You can see that there are two BGP groups with two neighbors each, making a total of four BGP neighbors.

```
[edit]
dhanks0@MX80# show policy-options prefix-list bgp-neighbors | display inheritance
##
## apply-path was expanded to:
## 10.0.3.3/32;
## 10.0.3.4/32;
## 10.0.9.6/32;
## 10.0.9.7/32;
##
apply-path "protocols bgp group <*> neighbor <*>";
```

```
[edit]
dhanks@MX80#
```

In this scenario, the `apply-path` begins at the `protocols bgp` stanza and matches all groups and all neighbors using the `<*>` wildcard. The last match catches the IP address and automatically fills out the prefix list `bgp-neighbors`.

ALERT! Using `apply-path` doesn't protect a prefix list against misconfiguration. If you add a bad BGP neighbor address the `apply-path` will automatically update the prefix list and allow the traffic to be passed though the firewall filter.

Apply-flags omit

Another method of reducing the Junos visual configuration is a hidden flag called *omit*. This option removes all child configuration from view and replaces it with `/* OMITTED */`. The `omit` keyword is useful for hiding large amounts of information that you won't use on a daily basis. Let's use an example by examining the details of the firewall filter `discard-all`:

```
dhanks@MX80# show firewall family inet filter discard-all term discard-tcp {
  from {
    protocol tcp;
  }
  then {
    count discard-tcp;
    log;
    discard;
  }
}
term discard-netbios {
```



```
    from {
      protocol udp;
      destination-port 137;
    }
    then {
      count discard-netbios;
      discard;
    }
  }
}
term discard-udp {
  from {
    protocol udp;
  }
  then {
    count discard-udp;
    log;
    discard;
  }
}
term discard-icmp {
  from {
    protocol icmp;
  }
  then {
    count discard-icmp;
    log;
    discard;
  }
}
term discard-unknown {
  then {
    count discard-unknown;
    log;
    discard;
  }
}
}
```

```
[edit]
dhanks@MX80#
```

This single firewall filter is nearly fifty lines in length and is generic enough that it is not going to be changed on a regular basis. It's a perfect candidate for being omitted while viewing your configuration using the `apply-flags omit` command:

```
[edit]
dhanks@MX80# set firewall family inet filter discard-all apply-flags omit
```

```
[edit]
dhanks@MX80#
```

Now, when your configuration is viewed with the `show` command, it is omitted:

```
family inet {
  filter discard-all { /* OMITTED */ };
```

```
[edit]
dhanks@MX80#
```

NOTE Because `omit` is a hidden command, you can't use tab-completion when configuring `apply-flags`. You need to type it out completely. It also will not show up in the inline help menu when you press `?`.

NOTE Configuration sections with the `apply-flags omit` are omitted only when viewed at a higher level. For example typing `show firewall family inet filter discard-all` would show the entire filter, whereas typing `show firewall family inet` shows it as being omitted.

Using the `apply-flags omit` has reduced this firewall filter from nearly fifty lines to a single line. But you should know that when you are viewing omitted configuration at a higher level, it's possible to still view the entire configuration using the command-line option `display omit`, as here:

```
[edit]
dhanks@MX80# show firewall family inet | display omit
filter discard-all {
  apply-flags omit;
  term discard-tcp {
    from {
      protocol tcp;
    }
    then {
      count discard-tcp;
      log;
      discard;
    }
  }
  term discard-netbios {
    from {
      protocol udp;
      destination-port 137;
    }
    then {
      count discard-netbios;
      discard;
    }
  }
}
```

```
}
term discard-udp {
  from {
    protocol udp;
  }
  then {
    count discard-udp;
    log;
    discard;
  }
}
term discard-icmp {
  from {
    protocol icmp;
  }
  then {
    count discard-icmp;
    log;
    discard;
  }
}
term discard-unknown {
  then {
    count discard-unknown;
    log;
    discard;
  }
}
}
```

Summary

The apply-path feature creates and maintains a prefix list based on the Junos configuration. The omit feature is verbose when you want it and terse when you need it. These simple tools pay big dividends. Let Junos work for you, not the other way around. Junos was designed to easily handle scale without creating operational burdens.

Chapter 5

Creating a Basic Framework of Firewall Filters

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Chapter 5 shows you how to create a basic framework of firewall filters using common routing protocols and routing engine services. These firewall filters will be the building blocks for creating a customized security profile to the routing engine. As always, it's recommended that you follow along on your test bed or device.

Overview of a Firewall Filter Framework

Throughout this chapter you will learn how to build a portable and customizable framework that can be used to secure the routing engine. There are three major components to this framework: generic policers, prefix-lists using apply-path, and firewall filters, as shown in Figure 5.1.

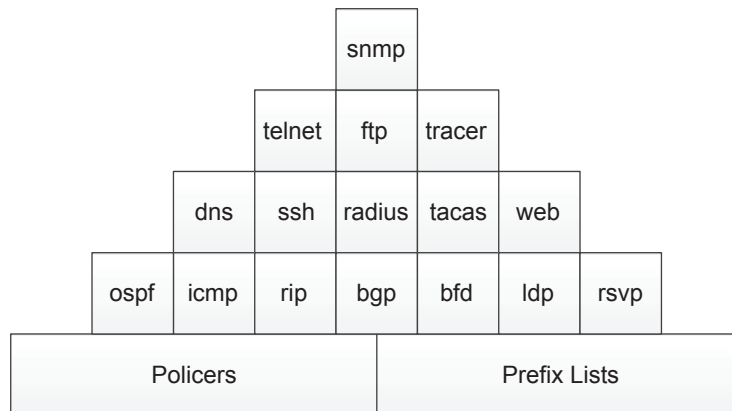


Figure 5.1 Framework for Securing the Routing Engine

You can see in Figure 5.1 that the policers and prefix lists are the foundation of the framework, with the firewall filters sitting on top and enlisting the policers and prefix lists. Each firewall filter is designed to only allow specific traffic, and the modular building blocks allow the administrator to cherry-pick firewall filters and apply them to the router, as shown in Figure 5.2.

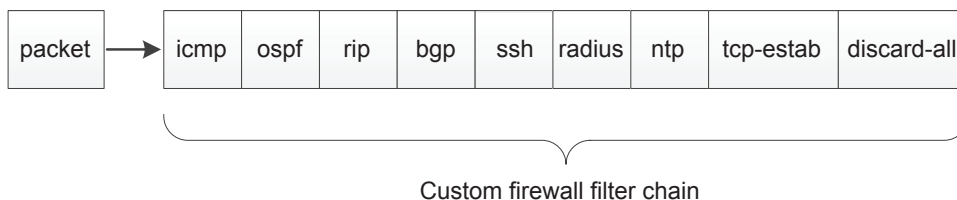


Figure 5.2 Example of Firewall Filter Chain Using “filter input-list”

The overall goal of this framework is that it should be portable enough to move into the majority of environments and “just work” without major modification. The prefix lists should be generic enough to include basic services, protocols, and the router’s IPv4 addresses. The firewall filters should be written so that each filter focuses on a specific problem, and when the prefix lists are built using `apply-path`, modification shouldn’t be necessary.

The firewall filters should be designed to be used on the control plane using the `interface 100.0` in the direction of input. The firewall filters are modular, so they’re able to be applied with the `input-list` option and be chained together to create custom firewall requirements for your particular router.

In short, this framework applies what was shown in the previous chapters of this book. You get modularity and customization at the same time, utilizing a single Junos operating system.

NOTE This chapter reviews and comments on a subset of the final framework configuration. The appendix contains the final working configuration in its entirety.

TIP A special Copy and Paste edition of this book can be downloaded from this book’s page at www.juniper.net/dayone. Its rich-text format allows you to copy and paste firewall filter configurations for your own use.

Prefix Lists

Prefix lists should use `apply-path` liberally when dealing with services, routing protocols, and router addresses, so that Junos is doing all of the redundant work and you can focus on more important things. The prefix lists can be referenced by the firewall filters without any modification as your network grows and changes

Local Addresses

Let’s start by creating prefix lists that define local addresses on the router. Depending on what type of hardware you are dealing with, there are different types of virtualization capabilities, so let’s use prefix lists that use IPv4 addresses defined on the router itself and inside of logical-systems.

Router IPv4 Addresses

Our prefix list `router-ipv4` contains all of the IPv4 addresses on the physical router, and the prefix list `router-ipv4-logical-systems` contains all of the IPv4 addresses that exist within any logical-system:

```
prefix-list router-ipv4 {
  apply-path "interfaces <*> unit <*> family inet address <*>";
}
prefix-list router-ipv4-logical-systems {
  apply-path "logical-systems <*> interfaces <*> unit <*> family inet address <*>";
}
```

NOTE The prefix-list `router-ipv4` is aggressive and will contain every single IPv4 address on the physical router, including GRE tunnels, loopback interfaces, and routing engine management interfaces such as `fxp0`.

Localhost

The loopback is usually defined as `127.0.0.1/32` in Junos, but the RFC 1700 defines it as `127/8`. Let's create a prefix list called `localhost` that matches `127/8`:

```
prefix-list localhost {
  127.0.0.0/8;
}
```

Protocols

Let's take the most common routing protocols and their multicast addresses and create easy to read prefix lists, so the configuration is a little more human-readable.

OSPF

OSPF uses two well-known multicast addresses to establish adjacency and send Hello packets:

IPv4 Multicast Address	Description
224.0.0.5	OSPF AllSPFRouters
224.0.0.6	OSPF AllDRouters

Let's create the prefix-list `ospf` to match both of these well-known multicast addresses:

```
prefix-list ospf {
  224.0.0.5/32;
  224.0.0.6/32;
}
```


RIP

RIPv2 uses a single well-known multicast address to exchange routing information, so let's create the prefix list `rip` to match this well-known multicast address:

```
prefix-list rip {
    224.0.0.9/32;
}
```

VRRP

Virtual Router Redundancy Protocol (VRRP) uses a single well-known multicast address that physical routers use to communicate, and our prefix list `vrrp` matches the multicast address 224.0.0.18:

```
prefix-list vrrp {
    224.0.0.18/32;
}
```

Multicast All Routers

The IPv4 multicast address 224.0.0.2 is reserved for AllRouters within a network segment. Protocols such as the Label Distribution Protocol (LDP) use this address to discover Hello packets. So let's create a prefix list called `multicast-all-routers` to match this well-known multicast address:

```
prefix-list multicast-all-routers {
    224.0.0.2/32;
}
```

BGP Neighbors

Let's use the `apply-path` feature to walk through the Junos configuration to discover all the BGP neighbors within the configuration.

```
prefix-list bgp-neighbors {
    apply-path "protocols bgp group <*> neighbor <*>";
}
prefix-list bgp-neighbors-logical-systems {
    apply-path "logical-systems <*> protocols bgp group <*>
neighbor <*>";
}
```

Services

The four most common services on a router are RADIUS, TACAS+, NTP, and SNMP. Let's build prefix lists to match the appropriate sections of the Junos configuration and find the server IP addresses.

RADIUS

Junos allows for multiple RADIUS servers to be defined under the `system radius-server` stanza. Let's create a prefix list to match all of the RADIUS servers:

```
system {
  radius-server {
    192.168.0.10 secret "$9$whYaZ"; ## SECRET-DATA
    192.168.0.11 secret "$9$5Q69"; ## SECRET-DATA
    192.168.0.12 secret "$9$DMimf"; ## SECRET-DATA
    192.168.0.13 secret "$9$ItYErE"; ## SECRET-DATA
  }
}
```

Even though the IP address is located in between `radius-server` and `secret`, you can still use wildcards to match the IP addresses. Let's create an `apply-path` that matches everything to the IP address and then stops, leaving out the secret:

```
prefix-list radius-servers {
  apply-path "system radius-server <*>";
}
```

This wildcard should match the four IP addresses 192.168.0.10 – 13, but's let's double-check that the `apply-path` is working as expected, by using the `display inheritance` command:

```
[edit]
dhanks@MX80# show policy-options prefix-list radius-servers | display inheritance
##
## apply-path was expanded to:
##   192.168.0.10;
##   192.168.0.11;
##   192.168.0.12;
##   192.168.0.13;
##
apply-path "system radius-server <*>";
```

```
[edit]
dhanks@MX80#
```

The double hash marks (`##`) verify that Junos was able to match the IP addresses and expand them out in a list.

TACAS+

TACAS+ is very similar to RADIUS in its configuration. Junos supports multiple TACAS+ servers defined in the `system tacplus-servers` stanza, as you can see here:

```

system {
  tacplus-server {
    172.16.0.100;
    172.16.0.101;
    172.16.0.102;
    172.16.0.103;
  }
}

```

Let's leverage the `apply-path` from the prefix list `radius-servers` and build a similar prefix list for TACAS+:

```

prefix-list tacas-servers {
  apply-path "system tacplus-server <*>";
}

```

Once again it should work, but let's verify the `apply-path` using the `display inheritance` command:

```

[edit]
dhanks@MX80 # show policy-options prefix-list tacas-servers | display inheritance
##
## apply-path was expanded to:
## 172.16.0.100/32;
## 172.16.0.101/32;
## 172.16.0.102/32;
## 172.16.0.103/32;
##
apply-path "system tacplus-server <*>";

```

```

[edit]
dhanks@MX80#

```

You can see the four TACAS+ servers 172.16.0.100-103 and the double hashes to verify that the `apply-path` is working as expected.

NTP

The Network Time Protocol (NTP) is defined under the `system ntp` stanza, and Junos supports multiple NTP servers:

```

system {
  ntp {
    server 192.168.33.10;
    server 192.168.33.11;
    server 192.168.33.12;
    server 192.168.33.13;
  }
}

```

Leveraging our experience with the previous `apply-path` prefix lists, let's apply the same methodology to NTP.

```
prefix-list ntp-servers {
    apply-path "system ntp server <*>";
}
```

And verify with `display inheritance`:

```
prefix-list ntp-server {
    ##
    ## apply-path was expanded to:
    ##   192.168.33.10;
    ##   192.168.33.11;
    ##   192.168.33.12;
    ##   192.168.33.13;
    ##
    apply-path "system ntp server <*>";
}
```

Here `apply-path` was able to successfully match the NTP servers and find the four IP addresses 192.168.33.10-13.

Junos also allows the use of NTP peers, so let's match those as well:

```
prefix-list ntp-server-peers {
    apply-path "system ntp peer <*>";
}
```

The greatest benefit to using `apply-path` is that Junos builds and maintains the dynamic prefix lists for you. As you add additional NTP servers to the configuration, there's no need to go back and modify the prefix list `ntp` because Junos does this automatically.

SNMP

The Simple Network Management Protocol (SNMP) supports multiple communities and the ability to define a different set of clients that are able to access each community. Junos enables the use of the `client-lists` keyword, which is similar to prefix lists.

ALERT! SNMP `client-list` and prefix lists *cannot* share the same name.

Example SNMP configuration:

```
snmp {
    client-list internal {
        10.0.0.0/8;
        192.168.0.0/16;
    }
}
```

```

        172.16.0.0/12;
    }
    community public {
        authorization read-only;
        clients {
            172.16.100.0/24;
        }
    }
    community private {
        authorization read-write;
        client-list-name internal;
    }
}

```

SNMP provides you with multiple levels of configuration and flexibility when defining which clients are allowed to access communities. Junos supports SNMP `client-list`, which acts as a prefix-list that can be applied to multiple SNMP communities.

To capture all of the SNMP clients, you need to define two prefix lists using `apply-path`. The first prefix list captures the statically defined clients and the other prefix list captures the SNMP `client-lists`.

```

prefix-list snmp-client-lists {
    apply-path "snmp client-list <*> <*>";
}
prefix-list snmp-community-clients {
    apply-path "snmp community <*> clients <*>";
}

```

Best practice is to always verify:

```

[edit]
dhanks@MX80# show policy-options prefix-list snmp-client-lists | display inheritance
##
## apply-path was expanded to:
## 10.0.0.0/8;
## 192.168.0.0/16;
## 172.16.0.0/12;
## 203.0.113.15/32;
## 203.0.113.16/32;
## 198.51.100.55/32;
## 198.51.100.56/32;
##
apply-path "snmp client-list <*> <*>";

```

```

[edit]
dhanks@MX80# show policy-options prefix-list snmp-community-clients | display
inheritance
##

```

```
## apply-path was expanded to:
## 172.16.100.0/24;
## 192.168.100.0/24;
##
apply-path "snmp community <*> clients <*>";
```

```
[edit]
dhanks@MX80#
```

You can see here that the prefix list `snmp-client-lists` successfully aggregated together all of the clients defined in the SNMP `client-lists`. The last prefix list `snmp-community-clients` aggregated all of the statically defined clients under each SNMP community. By using these two prefix lists together you're able to match any type of client combination within SNMP.

DNS

Junos allows for multiple DNS name servers to be defined under `system name-server`:

```
system {
  name-server {
    172.16.5.40;
    172.16.5.41;
    172.16.5.42;
    172.16.5.43;
  }
}
```

Let's create a prefix list to match the name servers defined:

```
prefix-list dns-servers {
  apply-path "system name-server <*>";
}
```

And when verified:

```
[edit]
dhanks@MX80# show policy-options prefix-list dns-servers | display inheritance
##
## apply-path was expanded to:
## 172.16.5.40/32;
## 172.16.5.41/32;
## 172.16.5.42/32;
## 172.16.5.43/32;
##
apply-path "system name-server <*>";
```

```
[edit]
dhanks@MX80#
```

Policers

Policing traffic on the routing engine is straightforward: define a bandwidth-limit and make the policer action discard any packets that exceed the bandwidth-limit. The burst-size-limit is the same across any policers applied to the routing engine because the link to the routing engine is 1Gbps, and the amount of bandwidth a 1 Gbps interface is able to send in 5ms is 625,000 bytes, which is the recommended value for the burst-size-limit.

Management

Let's create two generic management policers. These policers are term-specific so that a policer instance is created for each firewall filter term, and the policer actions will be to discard any packets that exceed the bandwidth limit.

The benefit of using two generic policers is that it reduces the size of the configuration and makes the configuration less complex, while still maintaining the firewall filter term policing. The alternative to using a generic policer is to define a policer per firewall filter or term. The benefit to using a term-specific policer is that when viewing the policer with the `show firewall` command, there's a policer instance for each firewall filter term. For example, if traffic for both SSH and SNMP had to be policed, it is possible to see how many packets of each were policed on a per firewall filter basis.

Previously it was recommended that firewall filter terms match the filter name and counter if possible. The value in such a naming convention now becomes apparent when trying to understand how many packets were policed for a particular firewall filter. Using this method allows you to create generic policers, but to apply them to specific firewall filters and maintain full visibility into policing statistics.

It's easily illustrated by viewing the output of `show firewall` and comparing the output between *term-specific* and *filter-specific* policers.

Term-specific (default)

Policers are term-specific by default. Each firewall filter term that has an action modifier of `policer` will have its own named policer instance.

Each policer is named after the firewall filter term from which it was applied. Because the firewall filter terms are named based on what we are filtering, it's easy to see what's being policed:

```
dhanks@MX80> show firewall filter lo0.3-i
```

```
Filter: lo0.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-bfd-lo0.3-i	3558932	68441
accept-bgp-lo0.3-i	40883	662
accept-dns-lo0.3-i	0	0
accept-icmp-lo0.3-i	0	0
accept-ldp-lo0.3-i	234426	3889
accept-ntp-lo0.3-i	0	0
accept-ospf-lo0.3-i	73096	1071
accept-rip-lo0.3-i	16640	320
accept-rip-igmp-lo0.3-i	32	1
accept-snmp-lo0.3-i	0	0
accept-ssh-lo0.3-i	0	0
accept-traceroute-lo0.3-i	0	0
accept-web-lo0.3-i	0	0
discard-icmp-lo0.3-i	0	0
discard-netbios-lo0.3-i	0	0
discard-tcp-lo0.3-i	0	0
discard-udp-lo0.3-i	0	0
discard-unknown-lo0.3-i	0	0
no-icmp-fragments-lo0.3-i	0	0

```
Policers:
```

Name	Packets
management-1m-accept-dns-lo0.3-i	0
management-1m-accept-ntp-lo0.3-i	0
management-1m-accept-traceroute-lo0.3-i	0
management-5m-accept-icmp-lo0.3-i	0
management-5m-accept-snmp-lo0.3-i	0
management-5m-accept-ssh-lo0.3-i	0
management-5m-accept-web-lo0.3-i	0

As you can see, it's easy to determine how many packets were policed for DNS, NTP, traceroute, ICMP, and other protocols. Because the policers are term-specific an instance of the policer is created for each term in the firewall filter and a counter is created in the naming format of `<policer_name>-<firewall_filter_term_name>-<interface_name>.<unit>-<direction>`.

Filter-specific

When the `filter-specific` knob is enabled on a policer and a firewall filter has multiple terms and action modifiers of `policer`, only a single policer instance will be created per firewall filter. This causes some loss

in visibility as to how many packets were policed on a per-filter or per-term basis. All of the terms within the firewall filter are policed as an aggregate.

```
dhanks@MX80> show firewall filter lo0.3-i
```

```
Filter: lo0.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-bfd-lo0.3-i	3464760	66630
accept-bgp-lo0.3-i	39776	644
accept-dns-lo0.3-i	0	0
accept-icmp-lo0.3-i	0	0
accept-ldp-lo0.3-i	228294	3787
accept-ntp-lo0.3-i	0	0
accept-ospf-lo0.3-i	71032	1041
accept-rip-lo0.3-i	16224	312
accept-rip-igmp-lo0.3-i	32	1
accept-snmp-lo0.3-i	0	0
accept-ssh-lo0.3-i	0	0
accept-traceroute-lo0.3-i	0	0
accept-web-lo0.3-i	0	0
discard-icmp-lo0.3-i	0	0
discard-netbios-lo0.3-i	0	0
discard-tcp-lo0.3-i	0	0
discard-udp-lo0.3-i	0	0
discard-unknown-lo0.3-i	0	0
no-icmp-fragments-lo0.3-i	0	0

```
Policers:
```

Name	Packets
management-1m-lo0.3-i	0
management-5m-lo0.3-i	0

It's impossible to calculate how many packets were policed on a per filter basis looking at these two policers. When policers are created with the `filter-specific` knob, a policer instance is created per firewall filter if there is an action modifier of `policer` and a counter is created. The naming convention is `<policer_name>-<interface_name>.<unit>-<direction>`.

BEST PRACTICE

Use term-specific policers (the default) when you need full visibility into what types of traffic are being policed in the routing engine.

Management 1 Mbps

The management policer `management-1m` will restrict traffic to 1 Mbps and discard any packets that exceed this bandwidth limit. This policer will be applied to protocols such as NTP, traceroute, RADIUS, TACAS+, and telnet. Traditionally these protocols do not require high throughput so they are a good candidate for this policer:

```

policer management-1m {
  apply-flags omit;
  if-exceeding {
    bandwidth-limit 1m;
    burst-size-limit 625k;
  }
  then discard;
}

```

Management 5 Mbps

The management policer `management-5m` is identical to the previous policer, except that traffic will have a bandwidth limit of 5 Mbps. Protocols such as ICMP, SNMP, HTTP, and SSH require higher throughput and 5 Mbps is adequate for accessing the router via SSH and copying configuration files via SCP:

```

policer management-5m {
  apply-flags omit;
  if-exceeding {
    bandwidth-limit 5m;
    burst-size-limit 625k;
  }
  then discard;
}

```

ALERT! It's recommended to remove *any policer* from SSH if you need to copy any large files to the routing engine such as a Junos firmware image.

Firewall Filters

Now it's time to begin creating the fundamental building blocks to securing the routing engine. Each firewall filter is specific to a routing protocol or service – in other words, no firewall filter restricts traffic *outside of its immediate focus*. Examine the framework gathered so far as illustrated in Figure 5.3.

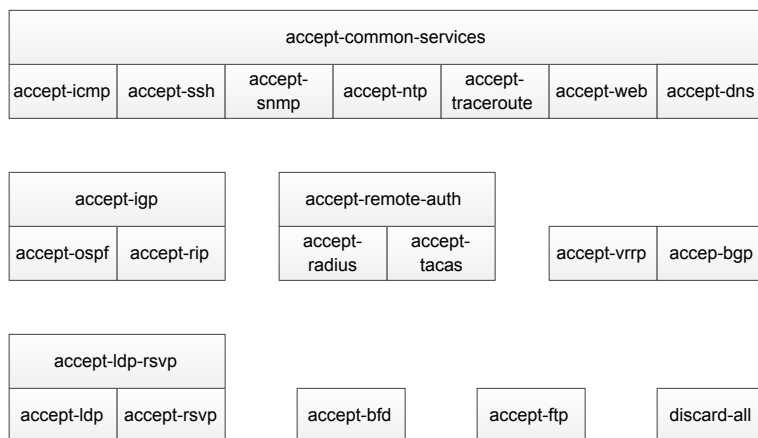


Figure 5.3 Firewall Filter Framework

You can see that firewall filters are grouped together where it makes sense. For example, the firewall filter `accept-igmp` includes both `accept-ospf` and `accept-rip`. This kind of grouping helps cut down on the number of firewall filters in the chain.

The firewall filters are designed to be used as a chain on the interface `lo0.0` to secure the routing engine. The order of the firewall filters in the chain isn't important except for the final "catch all" firewall filter `discard-all`.

It's important that the firewall filter `discard-all` be placed at the end of the chain for several reasons. Placing any firewall filters after `discard-all` causes all traffic to be denied and never evaluated by any firewall filters after `discard-all`. The `discard-all` firewall filter keeps track of how many packets were discarded, the total number of bytes discarded, and logs the packet's header.

Protocols

Let's now take a deep dive into all of the major protocols and how to create firewall filters around them. Our goal is to make the firewall filters specific enough to only allow certain types of traffic through, but generic enough so they should "just work" in most environments.

RIP

RIP is one of the easier routing protocols for which to write a firewall filter, because the destination-address is always `224.0.0.9`. Let's use the

prefix list `rip` that we created earlier. All the packets will be sourced from other routers, but within a *Direct* interface on the router. The prefix list `router-ipv4` contains all of the IPv4 addresses on the router and is used to match the source-address. Note that RIP also uses the protocol UDP which always uses the destination-port of 520/rip. Let's write the filter:

```
filter accept-rip {
  apply-flags omit;
  term accept-rip {
    from {
      source-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      destination-prefix-list {
        rip;
      }
      protocol udp;
      destination-port rip;
    }
    then {
      count accept-rip;
      accept;
    }
  }
  term accept-rip-igmp {
    from {
      source-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      destination-prefix-list {
        rip;
      }
      protocol igmp;
    }
    then {
      count accept-rip-igmp;
      accept;
    }
  }
}
```

You can also see a counter was added to each term so you can keep tabs on how many packets and bytes transferred have been passed through this filter. Good idea.

OSPF

OSPF uses two multicast addresses: 224.0.0.5 and 224.0.0.6. The multicast address 224.0.0.5 is used by OSPF to communicate with AllSPFRouters, while the address 224.0.0.6 is used to communicate with OSPF AllDRouters.

NOTE OSPF uses IP Protocol number 89/ospf. Since it uses its own protocol, you don't have to worry about port numbers or other attributes.

OSPF doesn't always use the destination-address of 224.0.0.5 or 224.0.0.6. The OSPF database description packets in specific interface types are sourced and destined between two routers using the interface IP addresses, so our filter should make sure to include both prefix lists `ospf` and `router-ipv4` as the destination-address to ensure that all the OSPF packet types are allowed through to the routing engine:

```
filter accept-ospf {
  apply-flags omit;
  term accept-ospf {
    from {
      source-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      destination-prefix-list {
        router-ipv4;
        ospf;
        router-ipv4-logical-systems;
      }
      protocol ospf;
    }
    then {
      count accept-ospf;
      accept;
    }
  }
}
```

BGP

BGP is deceptively easy, thanks to prefix lists using the `apply-path` feature, as the prefix list `bgp-neighbors` contains a list of all the BGP neighbors defined in protocols `bgp`.

BGP always uses the TCP protocol and either a destination or source port of 179/bgp. You can use the port `bgp` match condition to match either the destination or source port of BGP. The source address is

always the BGP neighbors as defined in protocols `bgp` so you can use the prefix list `bgp-neighbors`. The destination address is always an IP address on the router so you can use the prefix-list `router-ipv4` again:

```
filter accept-bgp {
  apply-flags omit;
  term accept-bgp {
    from {
      source-prefix-list {
        bgp-neighbors;
        bgp-neighbors-logical-systems;
      }
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      protocol tcp;
      port bgp;
    }
    then {
      count accept-bgp;
      accept;
    }
  }
}
```

VRRP

The VRRP protocol is simple. All VRRP packets are always sent to the multicast address 224.0.0.18. Let's leverage the prefix list `vrrp` we created earlier to make it easier to read.

VRRP uses the IP protocol 112/vrrp for all communication. The caveat is that if VRRP authentication is enabled, it uses Authentication Header (AH) which uses IP Protocol 51/ah. So our configuration needs to allow both protocols in this firewall filter for VRRP to function:

```
filter accept-vrrp {
  apply-flags omit;
  term accept-vrrp {
    from {
      source-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      destination-prefix-list {
        vrrp;
      }
      protocol [ vrrp ah ];
    }
    then {
      count accept-vrrp;
    }
  }
}
```

```

        accept;
    }
}

```

The `protocol [vrrp ah]` is a logical OR operation. Only one of the protocols in the list needs to match in order for `protocol` to be evaluated as true.

LDP

The Label Distribution Protocol (LDP) is straightforward and lends itself well to a firewall filter. Hellos are exchanged via UDP destined to the All Routers multicast address of 224.0.0.2 with a destination port of 646/ldp. Targeted LDP unicasts via UDP during the discovery phase instead of using multicast. After discovery LDP, it unicasts the peer directly via TCP destined to 224.0.0.2, with a destination port of 646/ldp.

This filter needs to create four terms to capture each specific phase:

- LDP discovery
- Targeted LDP discovery
- Unicast LDP
- IGMP join to 224.0.0.2

Let's review the filter together:

```

filter accept-ldp {
  apply-flags omit;
  term accept-ldp-discover {
    from {
      source-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      destination-prefix-list {
        multicast-all-routers;
      }
      protocol udp;
      destination-port ldp;
    }
    then {
      count accept-ldp-discover;
      accept;
    }
  }
  term accept-ldp-unicast {

```

```

    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        protocol tcp;
        port ldp;
    }
    then {
        count accept-ldp-unicast;
        accept;
    }
}
term accept-tldp-discover {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        protocol udp;
        destination-port ldp;
    }
    then {
        count accept-tldp-discover;
        accept;
    }
}
term accept-ldp-igmp {
    from {
        source-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        destination-prefix-list {
            multicast-all-routers;
        }
        protocol igmp;
    }
    then {
        count accept-ldp-igmp;
        accept;
    }
}
}
}

```

RSVP

The Resource Reservation Protocol (RSVP) is a transport layer protocol that uses IP protocol number 46. Traffic is always destined to the router so the filter can use the prefix list `router-ipv4`:


```
filter accept-rsvp {
  apply-flags omit;
  term accept-rsvp {
    from {
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      protocol rsvp;
    }
    then {
      count accept-rsvp;
      accept;
    }
  }
}
```

BFD

The Bidirectional Forwarding Protocol (BFD) uses UDP to transmit packets. There's a very specific range of ports that BFD is allowed to use for both the source and destination ports: the source ports allowed are 49152 through 65535 and the destination ports allowed are 3784 through 3785, which you'll see as part of the filter configuration:

```
filter accept-bfd {
  apply-flags omit;
  term accept-bfd {
    from {
      source-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      protocol udp;
      source-port 49152-65535;
      destination-port 3784-3785;
    }
    then {
      count accept-bfd;
      accept;
    }
  }
}
```

Common Services

The firewall filter *common services* groups commonly used protocols under a single firewall filter so that the `input-list` on the interface `lo0` doesn't become terribly long. Figure 5.4 illustrates the common services grouping concept.

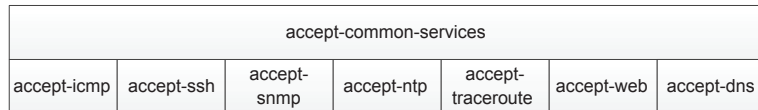


Figure 5.4 Common Services Concept

The firewall filter `accept-common-services` is designed to be used as the first firewall filter in the chain. It's constructed using the `filter` knob to reference another firewall filter, instead of reinventing the wheel and defining its own from and then statements:

```
filter accept-common-services {
  apply-flags omit;
  term accept-icmp {
    filter accept-icmp;
  }
  term accept-ssh {
    filter accept-ssh;
  }
  term accept-snm  
p {
    filter accept-snm  
p;
  }
  term accept-ntp {
    filter accept-ntp;
  }
  term accept-traceroute {
    filter accept-traceroute;
  }
  term accept-web {
    filter accept-web;
  }
  term accept-dns {
    filter accept-dns;
  }
}
```

The common services filter is designed so that you can modify it to fit your environment. Adding and removing filters has become a trivial task that can be completed with one command, for example, adding VRRP to `accept-common-services`:

```
dhanks@MX80# set firewall family inet filter accept-common-services term accept-vrrp
filter accept-vrrp
```

This command adds a term to the end of the `accept-common-services` filter called `accept-vrrp` and references the existing firewall filter `accept-vrrp`.

ICMP

The Internet Control Message Protocol (ICMP) is one of the core protocols in IP to signal errors and query messages. ICMP packets are very small in nature and there's no legitimate need to fragment these packets, so the first step in creating an ICMP filter is to match ICMP fragments, count and log them, and silently discard the packet.

Which ICMP types to allow into your network is a decision that you need to make. Some network operators like to be good network citizens and return meaningful error messages back to the sender, while other operators may prefer a more secure approach of not responding to errors or queries, keeping the sender in the dark.

Since ICMP is used primarily for errors and queries, by nature it shouldn't consume a lot of resources. Best practice is to rate-limit the number of ICMP packets that enter the router. Let's use the generic management policer of 5 Mbps:

```
filter accept-icmp {
  apply-flags omit;
  term no-icmp-fragments {
    from {
      is-fragment;
      protocol icmp;
    }
    then {
      count no-icmp-fragments;
      log;
      discard;
    }
  }
  term accept-icmp {
    from {
      protocol icmp;
      icmp-type [ echo-reply echo-request time-exceeded
unreachable source-quench router-advertisement parameter-
problem ];
    }
    then {
      policer management-5m;
      count accept-icmp;
      accept;
    }
  }
}
```

SSH

The best practice to access the router will always be SSH, which uses port 22/ssh and the TCP protocol. There are a couple of match conditions that need to be verified before applying this filter to a router.

ALERT! Double check that the match conditions in the firewall filter `accept-ssh` make sense for your environment before applying it to the control plane. This book's example assumes that the router will be accessed via SSH and will be accessed from a RFC1918 source address.

```
filter accept-ssh {
  apply-flags omit;
  term accept-ssh {
    from {
      source-prefix-list {
        rfc1918;
      }
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      protocol tcp;
      port ssh;
    }
    then {
      policer management-5m;
      count accept-icmp;
      accept;
    }
  }
}
```

It is best practice to rate limit SSH to a reasonable bandwidth limit, usually a value of 5 Mbps. This is more than adequate for any CLI access for multiple users. The only drawback is that it takes about five or six minutes to SCP a version of Junos to the router. If this is unacceptable during a maintenance window, it's easy enough to temporarily deactivate the policer, like so:

```
[edit]
dhanks@MX80# deactivate firewall family inet filter accept-ssh term accept-ssh then
policer
```

To activate the policer after the upgrade is complete use the `activate` command:

```
[edit]
dhanks@MX80# activate firewall family inet filter accept-ssh term accept-ssh then
policer
```

TIP Don't forget that after activating or deactivating a section of the Junos configuration, you still need to commit the changes.

SNMP

The SNMP protocol uses the UDP protocol on port 161/snmp. Care needs to be taken to review the requirements of SNMP. The author has assumed that network management software accessing the router will have RFC1918 source addresses.

The bandwidth required for SNMP polling isn't very high. Best practice is to apply a policer to SNMP traffic. We'll use a bandwidth limit of 5 Mbps.

The source address can be gathered from the Junos configuration of SNMP via the statically defined community clients or SNMP `client-lists`. You can use the prefix list `snmp-client-lists` and `snmp-community-clients`. The destination address will always be an interface on the router, so you can use the prefix list `router-ipv4`, like so:

```
filter accept-snmp {
  apply-flags omit;
  term accept-snmp {
    from {
      source-prefix-list {
        snmp-client-lists;
        snmp-community-clients;
      }
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      protocol udp;
      destination-port snmp;
    }
    then {
      policer management-5m;
      count accept-snmp;
      accept;
    }
  }
}
```

NTP

The Network Time Protocol (NTP) is able to operate in several different modes: client, server, and symmetric active (peer) mode. Let's break up the NTP firewall filter into three terms so that you're able to differentiate the various types of NTP traffic:

```
filter accept-ntp {
  apply-flags omit;
  term accept-ntp {
    from {
      source-prefix-list {
        ntp-server;
        localhost;
      }
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
        localhost;
      }
      protocol udp;
      port ntp;
    }
    then {
      policer management-1m;
      count accept-ntp;
      accept;
    }
  }
}
term accept-ntp-peer {
  from {
    source-prefix-list {
      ntp-server-peers;
    }
    destination-prefix-list {
      router-ipv4;
      router-ipv4-logical-systems;
    }
    protocol udp;
    destination-port ntp;
  }
  then {
    policer management-1m;
    count accept-ntp-peer;
    accept;
  }
}
term accept-ntp-server {
  from {
    source-prefix-list {
      rfc1918;
    }
    destination-prefix-list {
      router-ipv4;
      router-ipv4-logical-systems;
    }
    protocol udp;
    destination-port ntp;
  }
}
```

```

        then {
            policer management-1m;
            count accept-ntp-server;
            accept;
        }
    }
}

```

Traceroute

Traceroute is a tool that varies wildly in implementation. Most networking equipment and UNIX variants default to using UDP with destination ports between 33435 and 33450. Other versions of traceroute, such as the one for Windows, use the ICMP protocol instead. Even when using the ICMP version of traceroute there are variations in the ICMP type. Newer versions of traceroute, such as “tcptraceroute”, use TCP and allow the user to specify a destination port. Tools like this come in handy if there is a firewall in the path of the traceroute.

Let’s create three terms to identify all of the major variations of traceroute and account for them.

NOTE It may seem like common sense to add a final term to discard all unknown traceroute traffic, but this is difficult to do. Protocols such as RSVP, OSPF, and RIP use packets with a time to live (TTL) of 1. Later in this chapter the final firewall filter “discard-all” is reviewed and this is a good candidate to catch and discard unknown traceroute traffic.

Here’s the traceroute configuration:

```

filter accept-traceroute {
    apply-flags omit;
    term accept-traceroute-udp {
        from {
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol udp;
            ttl 1;
            destination-port 33435-33450;
        }
        then {
            policer management-1m;
            count accept-traceroute-udp;
            accept;
        }
    }
}

```

```

term accept-traceroute-icmp {
  from {
    destination-prefix-list {
      router-ipv4;
      router-ipv4-logical-systems;
    }
    protocol icmp;
    ttl 1;
    icmp-type [ echo-request timestamp time-exceeded ];
  }
  then {
    policer management-1m;
    count accept-traceroute-icmp;
    accept;
  }
}
term accept-traceroute-tcp {
  from {
    destination-prefix-list {
      router-ipv4;
      router-ipv4-logical-systems;
    }
    protocol tcp;
    ttl 1;
  }
  then {
    policer management-1m;
    count accept-traceroute-tcp;
    accept;
  }
}
}

```

All traffic identified as traceroute is matched and accounted for. The traceroute packets are also policed down to 1 Mbps so that the routing engine isn't wasting CPU cycles responding to low-priority traffic.

Web

Junos allows for different types of management access. To allow J-Web with a firewall filter, you just need to match TCP packets that are destined to ports 80/http or 443/https. Let's police all J-Web traffic to 5 Mbps, as the service isn't that bandwidth intensive.

```

filter accept-web {
  apply-flags omit;
  term accept-web {
    from {
      source-prefix-list {

```



```

        rfc1918;
    }
    destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
    }
    protocol tcp;
    destination-port [ http https ];
}
then {
    policer management-5m;
    count accept-web;
    accept;
}
}
}
}

```

DNS

Junos uses the Domain Name System (DNS) to lookup hostnames from IP addresses. This traffic is initiated from Junos to the external DNS server, so let's create a firewall filter to match the return traffic. The DNS server replies back with a source port of 53/dns. Let's leverage the prefix list `dns-servers` to match the DNS name servers already defined in the Junos configuration:

```

filter accept-dns {
    apply-flags omit;
    term accept-dns {
        from {
            source-prefix-list {
                dns-servers;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol udp;
            source-port 53;
        }
        then {
            policer management-1m;
            count accept-dns;
            accept;
        }
    }
}
}

```

DNS queries are very small packets and shouldn't use a lot of bandwidth, so install a 1 Mbps policer for all DNS traffic.

Less Common Services

Two of the less common services used are FTP and telnet. These are insecure protocols as they transmit data in clear text. These services weren't included in the firewall filter `accept-common-services` because there are better alternatives. Instead of using FTP or telnet, it's much easier to use SSH. SSH uses the Diffie-Hellman key exchange and provides a secure transport for both command-line access and copying files.

ALERT! Do not use the following firewall filters unless absolutely necessary. FTP and telnet are insecure and pass data in clear text. Please use SSH and SCP instead.

FTP

The File Transfer Protocol (FTP) has two modes of operation: passive and active. The primary difference is how the data port is negotiated and established. With active mode, the FTP client connects from a random unprivileged port (N) to the FTP server's command port (21). The FTP client then listens on data port N+1. The FTP server then initiates a connection to the FTP client on port N+1.

Creating a stateless firewall filter for active FTP is straightforward. FTP uses the TCP protocol and the ports are always 20/ftp-data and 21/ftp, depending on whether the router is acting as a FTP server or client.

Firewall filters are stateless by default. It's out of the scope of the book to create a stateful firewall filter.

ALERT! This firewall filter isn't required to allow Junos to be an FTP client. Later in this chapter, the firewall filter `accept-tcp-established` will allow Junos to act as a FTP client.

```
filter accept-ftp {
  apply-flags omit;
  term accept-ftp {
    from {
      source-prefix-list {
        rfc1918;
      }
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
    }
    protocol tcp;
  }
}
```

```

        port [ ftp ftp-data ];
    }
    then {
        policer management-5m;
        count accept-ftp;
        accept;
    }
}

```

Telnet

The author was reluctant to include this filter, but understands there's always a case for its use. Telnet uses the TCP protocol and traffic is always destined to port 23/telnet.

ALERT! Don't use the firewall filter `accept-telnet` unless you know what you're doing.

```

filter accept-telnet {
    apply-flags omit;
    term accept-telnet {
        from {
            source-prefix-list {
                rfc1918;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol tcp;
            destination-port telnet;
        }
        then {
            policer management-1m;
            count accept-telnet;
            accept;
        }
    }
}

```

Telnet isn't a bandwidth intensive protocol so you should police it down to 1 Mbps.

Catch All

There is now a large framework of firewall filters, policers, and prefix lists to allow specified traffic into the routing engine such as OSPF, BGP, and SSH. Yet there will always be traffic that will ultimately not

match any of the filters applied to the routing engine. So to gain visibility into the traffic that is destined to the routing engine, but not explicitly allowed by the framework, there are three *catch all* firewall filters that can be applied at the very end of the firewall filter chain: TCP Established, Discard All, and Accept All.

TCP Established

It's important to note that these firewall filters have been designed to be applied to the ingress traffic on the routing engine. Traffic egressing the routing engine such as SSH, FTP, fetch, and TFTP is assumed to be trusted.

To match the return traffic that the routing engine initiated, you need to create a firewall filter that accepts TCP traffic that's considered `tcp-established`, which is synonymous with TCP flags matching ACK or RST.

You could simply create a firewall filter that allowed any TCP packets with the ACK or RST in the TCP flags, but it would be better to understand which protocols and services the packets belong to.

To gain visibility into the different protocols, the firewall filter breaks `accept-established` into seven terms:

- **SSH:** The term `accept-tcp-established-ssh` will match SSH return traffic back into the routing engine. The protocol will always be TCP with a source port of 22.
- **FTP:** The term `accept-tcp-established-ftp` will accept return traffic from the FTP command port. This will always be the TCP protocol with a source port of 21/ftp.
- **FTP DATA TCP-INITIAL:** Before the file transfer begins, there needs to be a negotiation of which FTP data port to use. The FTP server will initiate a connection to the router. This will always use the TCP protocol with the TCP flags SYN without an ACK. The synonym for this TCP flag is `tcp-initial`. The source port will always be 20/ftp-data.
- **FTP DATA:** After the FTP server has opened a data port to the router, data can be transferred. This traffic will always use the TCP protocol, have the TCP flags ACK or RST (`tcp-established`), and have the source port of 20/ftp-data.
- **TELNET:** The telnet protocol always uses TCP and when used from the command-line to connect to a remote host, the return

traffic will always have the source port of 23/telnet and the tcp-established TCP flags.

- **FETCH:** Junos supports a `fetch` command in the shell that allows the administrator to download HTTP or HTTPS content to the routing engine. Junos also supports CLI options to remotely download configurations with the `load` command while in configuration mode.
- The HTTP or HTTPS traffic will always use the TCP protocol, have a source port of either 80/http or 443/https, and have the tcp-established TCP flags.
- **EPHEMERAL:** Protocols such as the Trivial File Transfer Protocol (TFTP) will use UDP with a destination port of 69/tftp. An important note is that the server and client will negotiate ephemeral ports to be used for the data transfer. IANA suggests ports 49152 through 65535 to be used as dynamic and/or private ports.

ALERT! The author has found that the ephemeral port range of 49152-65535 works well for TFTP, but operating systems such as Linux and Windows can use different ranges of ephemeral ports. If you need to support additional protocols or operating systems that support a different range of ephemeral ports, you need to change the values in the term `accept-ephemeral`.

And here's the filter configuration:

```
filter accept-established {
  term accept-established-tcp-ssh {
    from {
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      source-port ssh;
      tcp-established;
    }
    then {
      policer management-5m;
      count accept-established-tcp-ssh;
      accept;
    }
  }
  term accept-established-tcp-ftp {
    from {
      destination-prefix-list {
        router-ipv4;
```

```
        router-ipv4-logical-systems;
    }
    source-port ftp;
    tcp-established;
}
then {
    policer management-5m;
    count accept-established-tcp-ftp;
    accept;
}
}
term accept-established-tcp-ftp-data-syn {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        source-port ftp-data;
        tcp-initial;
    }
    then {
        policer management-5m;
        count accept-established-tcp-ftp-data-syn;
        accept;
    }
}
term accept-established-tcp-ftp-data {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        source-port ftp-data;
        tcp-established;
    }
    then {
        policer management-5m;
        count accept-established-tcp-ftp-data;
        accept;
    }
}
term accept-established-tcp-telnet {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        source-port telnet;
        tcp-established;
    }
    then {
        policer management-5m;
```

```

        count accept-established-tcp-telnet;
        accept;
    }
}
term accept-established-tcp-fetch {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        source-port [ http https ];
        tcp-established;
    }
    then {
        policer management-5m;
        count accept-established-tcp-fetch;
        accept;
    }
}
term accept-established-udp-ephemeral {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        protocol udp;
        destination-port 49152-65535;
    }
    then {
        policer management-5m;
        count accept-established-udp-ephemeral;
        accept;
    }
}
}
}

```

It may seem easier to create a filter that simply accepts all TCP traffic with the `tcp-established` flag. The drawback to this method is that all return traffic would be aggregated and the network operator wouldn't be able to gain visibility into the traffic.

Discard All

The last firewall filter of the catch all is `discard-all`. The purpose of this filter is to simply partition packets into protocol categories that the routing engine isn't interested in and discard the packets silently.

The firewall filter `discard-all` should be the very last filter in the firewall filter chain applied to the `input-list` on the loopback interface.

ALERT! If the firewall filter `discard-all` isn't applied at the very end of the `input-list` chain, you will inadvertently drop all traffic destined to the routing engine without visibility.

The filter `discard-all` is the last line of defense to the routing engine. It's important that this filter arm the administrator with detailed statistics regarding what packets were denied.

The filter breaks `discard-all` into seven terms: `discard-ip-options`; `discard-TTL_1-unknown`; `discard-tcp`; `discard-netbios`; `discard-udp`; `discard-icmp`; and finally `discard-unknown`. Each term has a counter named after the term so that you know how many bytes and packets have been discarded. The packet header will also be logged twice. The first log is the PFE, allowing you to quickly check and see detailed packet header information with the `show firewall log` command. The second log is exported to the syslog servers defined in `system syslog host`, allowing storage and archival of the offending traffic that can be analyzed and reported on by tools such as Juniper's Security Threat Response Manager (STRM).

MORE? For more information about Juniper STRM see <http://www.juniper.net/us/en/products-services/security/strm-series/>.

And here is the `discard-all` filter's configuration:

```
filter discard-all {
  apply-flags omit;
  term discard-ip-options {
    from {
      ip-options any;
    }
    then {
      count discard-ip-options;
      log;
      syslog;
      discard;
    }
  }
  term discard-TTL_1-unknown {
    from {
      ttl 1;
    }
    then {
      count discard-all-TTL_1-unknown;
      log;
      syslog;
      discard;
    }
  }
}
```



```
term discard-tcp {
  from {
    protocol tcp;
  }
  then {
    count discard-tcp;
    log;
    syslog;
    discard;
  }
}
term discard-netbios {
  from {
    protocol udp;
    destination-port 137;
  }
  then {
    count discard-netbios;
    log;
    syslog;
    discard;
  }
}
term discard-udp {
  from {
    protocol udp;
  }
  then {
    count discard-udp;
    log;
    syslog;
    discard;
  }
}
term discard-icmp {
  from {
    protocol icmp;
  }
  then {
    count discard-icmp;
    log;
    syslog;
    discard;
  }
}
term discard-unknown {
  then {
    count discard-unknown;
    log;
    syslog;
    discard;
  }
}
```

```
    }
}
```

To demonstrate the power of the `discard-all` firewall filter, let's send some traffic to the routing engine that will be discarded, and then view how many packets and bytes were discarded, as well as any detailed packet header information.

Let's use a web browser and point it at the router via the URL `http://172.16.1.11:666`. Since there's no firewall filter allowing TCP traffic destined to port 666, it should be caught by the catch all firewall filter `discard-all`.

Since the term `discard-tcp` in the filter has an action modifier of `syslog`, let's take a look at the remote syslog server and view the packet header information:

```
Apr 02 22:43:13 172.16.1.11 Apr  3 05:42:49 /kernel: FW: fxp0.0      D tcp 172.16.1.100
172.16.1.11 53232 666
Apr 02 22:43:17 172.16.1.11 Apr  3 05:42:52 /kernel: FW: fxp0.0      D tcp 172.16.1.100
172.16.1.11 53232 666
Apr 02 22:43:24 172.16.1.11 Apr  3 05:42:58 /kernel: FW: fxp0.0      D tcp 172.16.1.100
172.16.1.11 53232 666
```

You could also confirm the same information by viewing the firewall log:

```
dhanks@MX80> show firewall log
Log :
Time      Filter  Action Interface  Protocol  Src Addr  Dest Addr
05:42:58  lo0.0-i  D    fxp0.0      TCP       172.16.1.100 172.16.1.11
05:42:52  lo0.0-i  D    fxp0.0      TCP       172.16.1.100 172.16.1.11
05:42:49  lo0.0-i  D    fxp0.0      TCP       172.16.1.100 172.16.1.11
```

You can also see the three packets were found by the counter `discard-tcp-lo0.0-i`:

```
dhanks@MX80> show firewall filter lo0.0-i counter discard-tcp-lo0.0-i

Filter: lo0.0-i
Counters:
Name                               Bytes          Packets
discard-tcp-lo0.0-i                152            3

dhanks@MX80>
```

Accept All

The firewall filter `accept-all` is the exact opposite of `discard-all`. It exists as an alternative to `discard-all` if the administrator is in a

production environment and needs to err on the side of caution when applying firewall filters to the routing engine. Use the filter `accept-all` in place of `discard-all` at the very end of the firewall filter chain.

NOTE The firewall filter `accept-all` is a safe alternative to `discard-all` as it will accept, log, and count all traffic. Only use `accept-all` if you're in a production environment and need to verify that the previous firewall filters in the chain are working as expected. The filter `accept-all` exists so that the administrator can see any traffic that was not explicitly accepted by previous filters in the chain and make corrections if necessary.

The filter is broken down by protocols and has an unknown catch-all at the very end to ensure that all traffic is accepted.

ALERT! Once you verify that all traffic is being matched by firewall filters before the `accept-all`, it's recommended to immediately remove `accept-all` and replace it with `discard-all` since the entire point is to secure the routing engine.

And here is the `accept-all` configuration:

```
filter accept-all {
  apply-flags omit;
  term accept-all-tcp {
    from {
      protocol tcp;
    }
    then {
      count accept-all-tcp;
      log;
      syslog;
      accept;
    }
  }
  term accept-all-udp {
    from {
      protocol udp;
    }
    then {
      count accept-all-udp;
      log;
      syslog;
      accept;
    }
  }
  term accept-all-igmp {
    from {
```

```

        protocol igmp;
    }
    then {
        count accept-all-igmp;
        log;
        syslog;
        accept;
    }
}
term accept-icmp {
    from {
        protocol icmp;
    }
    then {
        count accept-all-icmp;
        log;
        syslog;
        accept;
    }
}
term accept-all-unknown {
    then {
        count accept-all-unknown;
        log;
        syslog;
        accept;
    }
}
}
}

```

Each type of protocol is separated into different terms and counters. All traffic is logged to both the PFE and syslog for analysis. Any traffic that can't be identified will be accepted by the last term `accept-all-unknown`.

Summary

This chapter has demonstrated how to create a framework of firewall filters that can be used to secure various types of routers. The filters have been designed to be used in a firewall filter chain. The next-to-last filter should be the filter `accept-established`. When traffic is sourced from the routing engine, return traffic needs to be let back in. The last filter in the chain should be one of the two “catch all” filters: `accept-all` or `discard-all`.

Chapter 6

Applying Security Policies to the Routing Engine

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Building a framework to secure the routing engine that's portable, simple, yet customized to meet the requirements of the majority of networks has required the integration of a lot of moving parts. Here's a quick summary of what you've accomplished so far:

- *Chapter 1: Firewall Filters:* This chapter explained firewall filters in detail – how they match traffic, perform actions, and how they're evaluated. You also reviewed intermediate to expert level topics such as nested firewall filter, and firewall chaining.
- *Chapter 2: Policers:* This chapter introduced policers and how they rate limit traffic. You reviewed the Token Bucket algorithm in detail so that we could configure `bandwidth-limit` and `burst-size-limit` with confidence. You closed the chapter with how policer instances are created and how to influence the number of policers per firewall filter.
- *Chapter 3: Viewing Counters, Logs, and Policers:* This chapter demonstrated how to view counters, logs, and policers. You reviewed the Junos counter and policer naming convention.
- *Chapter 4: Junos Configuration Automation:* This chapter introduced a powerful Junos configuration automation tool called `apply-path`. You created example configuration and generated dynamic prefix lists based on the content of the Junos configuration.
- *Chapter 5: Creating a Basic Framework of Firewall Filters:* This chapter built upon the previous chapters and introduced the concept of a framework that's capable of providing building blocks that can be used in any combination to provide custom security to the routing engine. Each firewall filter within the framework was reviewed in detail.

Now let's take all of the material you've learned in Chapters 1 through 5 and apply a security policy to our router.

This book includes a complete example configuration located in the appendix, and you'll start by taking this configuration and loading it into your router. Then you will test a very simple security framework on a single router and gradually move up to a multiple router environment running multiple protocols and services.

TIP

If you are following along on your router, a special Copy and Paste edition of this book can be downloaded from this book's page at www.

juniper.net/dayone. Its rich-text format allows you to copy and paste the firewall filter configurations without having to key them in again.

Before You Begin

Adding firewall filters to the routing engine shouldn't be taken lightly. Any mistake could leave your router isolated and possibly require a connection to the console port to remove the offending firewall filter. This book calls out these specific pitfalls and warnings along the way.

Applying the configuration for the first time, however, needs to be done on a lab router, as this book has advocated all along. This way, if there are any mistakes they will be isolated to the laboratory equipment and not impact your production network.

Load Configuration

Let's start off by loading the framework configuration.

Please reference the configuration listed in the Appendix: *Framework Configuration*. The configuration includes prefix lists, firewall filters, and policers.

To Load the Configuration

1. Log in to the router and enter configuration mode.

```
dhanks@MX80> configure  
Entering configuration mode
```

```
[edit]  
dhanks@MX80#
```

2. Copy (CTRL-C) the entire configuration as shown in the Appendix, Framework Configuration. Use the `load merge relative terminal` command to import the configuration into your router.

```
dhanks@MX80# load merge relative terminal  
[Type ^D at a new line to end input]
```

3. Paste (CTRL-V) the configuration into the terminal. Press *enter* a couple of times after the entire configuration has been pasted into the terminal. Now type CTRL-D to end the load command.

```
<paste framework configuration>  
CTRL-D  
load complete
```

```
[edit]
dhanks@MX80#
```

4. Use the `show firewall` command to verify that the configuration has been imported correctly. The firewall filters should appear exactly as below with the `apply-flags omit`.

```
[edit]
dhanks@MX80# show firewall
family inet {
  filter accept-bgp { /* OMITTED */ };
  filter accept-ospf { /* OMITTED */ };
  filter accept-rip { /* OMITTED */ };
  filter accept-vrrp { /* OMITTED */ };
  filter accept-icmp { /* OMITTED */ };
  filter accept-ssh { /* OMITTED */ };
  filter accept-snmp { /* OMITTED */ };
  filter accept-ntp { /* OMITTED */ };
  filter accept-web { /* OMITTED */ };
  filter discard-all { /* OMITTED */ };
  filter accept-traceroute { /* OMITTED */ };
  filter accept-igp { /* OMITTED */ };
  filter accept-common-services { /* OMITTED */ };
  filter accept-bfd { /* OMITTED */ };
  filter accept-ldp { /* OMITTED */ };
  filter accept-ftp { /* OMITTED */ };
  filter accept-rsvp { /* OMITTED */ };
  filter accept-radius { /* OMITTED */ };
  filter accept-tacas { /* OMITTED */ };
  filter accept-remote-auth { /* OMITTED */ };
  filter accept-telnet { /* OMITTED */ };
  filter accept-dns { /* OMITTED */ };
  filter accept-ldp-rsvp { /* OMITTED */ };
  filter accept-established { /* OMITTED */ };
}
policer management-1m { /* OMITTED */ };
policer management-5m { /* OMITTED */ };
```

```
[edit]
dhanks@MX80#
```

5. Delete any existing firewall filters from your `lo0.0` interface.

```
[edit]
dhanks@MX80# delete interfaces lo0.0 family inet filter
```

6. Use the `commit check` command to ensure that the new configuration is syntactically correct and there are no errors.

```
dhanks@MX80# commit check
configuration check succeeds
```


7. Commit the configuration now that it has cleared the check.

```
dhanks@MX80# commit
commit complete
[edit]
dhanks@MX80#
```

NOTE If there were any errors during the load command, make sure that you correctly copied the entire configuration and pasted the entire configuration into your terminal.

TIP If there are still load errors, it may be because there are existing prefix lists, firewalls, or policers that are causing a naming conflict.

At this point your router should have a fully working framework for applying security to the routing engine. At the moment nothing is being secured because we deleted any existing filters applied to the interface `lo0.0`.

Your First Security Policy

Our first security policy will be very simple and only deal with one router. You'll apply the firewall filter `accept-common-services` and also include both "catch all" filters `accept-established` and `discard-all`.

The first filter includes the following services: ICMP, traceroute, SSH, SNMP, NTP, web, and DNS.

To Build the First Security Policy

1. Apply the firewall filter `accept-common-services` as the first filter in the `input-list` chain on the interface `lo0.0`.

```
[edit]
dhanks@MX80# set interfaces lo0.0 family inet filter input-list accept-common-services
```

2. Apply the first catch all firewall filter `accept-established`.

```
[edit]
dhanks@MX80# set interfaces lo0.0 family inet filter input-list accept-established
```

3. Apply the second catch all firewall filter `discard-all`.

```
[edit]
dhanks@MX80# set interfaces lo0.0 family inet filter input-list discard-all
```

4. Verify the changes with `show compare`.

```
dhanks@MX80# show | compare
```

```
[edit interfaces lo0 unit 0 family inet]
+   filter {
+     input-list [ accept-common-services accept-established discard-all ];
+   }
```

```
[edit]
dhanks@MX80#
```

5. Commit the changes using the `confirmed` feature. If there are any problems after the commit, the configuration will automatically rollback after 5 minutes.

```
dhanks@MX80# commit confirmed 5
commit confirmed will be automatically rolled back in 5 minutes unless confirmed
commit complete
```

```
# commit confirmed will be rolled back in 5 minutes
[edit]
dhanks@MX80#
```

6. After verifying that you still have connectivity to the router, issue the `commit` command again to confirm the configuration.

```
[edit]
dhanks@MX80# commit
commit complete
```

You have now applied your first security framework to the router. Let's take a look and see what's happening under the hood. What traffic is being allowed? Is there any traffic being denied? Take what you learned in Chapter 3 and begin reviewing the firewall counters, policers, and logs.

Let's take a look at the firewall filters applied to the routing engine by using the `show firewall filter` command:

```
dhanks@MX80> show firewall filter lo0.0-i
```

```
Filter: lo0.0-i
```

```
Counters:
```

Name	Bytes	Packets
accept-established-tcp-fetch-lo0.0-i	0	0
accept-established-tcp-ftp-lo0.0-i	0	0
accept-established-tcp-ftp-data-lo0.0-i	0	0
accept-established-tcp-ftp-data-syn-lo0.0-i	0	0
accept-established-tcp-ssh-lo0.0-i	0	0
accept-established-tcp-telnet-lo0.0-i	0	0
accept-established-udp-ephemeral-lo0.0-i	0	0
accept-icmp-lo0.0-i	0	0
accept-ntp-lo0.0-i	0	0
accept-ntp-server-lo0.0-i	0	0
accept-snmp-lo0.0-i	0	0
accept-ssh-lo0.0-i	3848	52

```

accept-traceroute-icmp-100.0-i          0      0
accept-traceroute-tcp-100.0-i          0      0
accept-traceroute-udp-100.0-i          0      0
accept-web-100.0-i                      0      0
discard-all-TTL_1-unknown-100.0-i     0      0
discard-icmp-100.0-i                   0      0
discard-netbios-100.0-i                 234    3
discard-tcp-100.0-i                    0      0
discard-udp-100.0-i                    0      0
discard-unknown-100.0-i                0      0
no-icmp-fragments-100.0-i              0      0
Policers:
Name                                     Packets
management-1m-accept-ntp-100.0-i       0
management-1m-accept-ntp-server-100.0-i 0
management-1m-accept-traceroute-icmp-100.0-i 0
management-1m-accept-traceroute-tcp-100.0-i 0
management-1m-accept-traceroute-udp-100.0-i 0
management-5m-accept-established-tcp-fetch-100.0-I 0
management-5m-accept-established-tcp-ftp-100.0-i 0
management-5m-accept-established-tcp-ftp-data-100.0-i 0
management-5m-accept-established-tcp-ftp-data-syn-100.0-i 0
management-5m-accept-established-tcp-ssh-100.0-i 0
management-5m-accept-established-tcp-telnet-100.0-i 0
management-5m-accept-established-udp-ephemeral-100.0-i 0
management-5m-accept-icmp-100.0-i      0
management-5m-accept-snmp-100.0-i      0
management-5m-accept-ssh-100.0-i       0
management-5m-accept-web-100.0-i       0

```

dhanks@MX80>

There's not too much going on here in this example of output. You can see that it's already counting the number of bytes and packets for SSH by the author being logged into the router via SSH.

NOTE If you happen to be savvy, the author's PC is on the same network as the routing engine causing some NetBIOS traffic to be discarded.

Let's try generating some traffic so you can see the counters and policers kick in. Try some of the following on your lab router:

- Using your PC or another router, use the traceroute command to the router.
- Using your PC or another router, use the SCP command and copy a file to the router.
- Using your PC or another router, use the ping command to the router.

- Using your PC or another router, use the ping command, but change the packet size to 9000. This will cause ICMP to fragment the packet (assuming your PC isn't using Jumbo frames).
- Using the router, telnet to another router or device.
- Using the router, FTP to another router or device and copy a file.

This should generate some interesting traffic and increase the counters, causing some traffic to be policed and some discarded in the process. Let's take a look at the counters and see what has changed and how the router responded to the traffic. Your lab router will be slightly different, of course. Use the `show firewall filter` command to view the counters and policers on interface `lo0.0`:

```
dhanks@MX80> show firewall filter lo0.0-i
```

```
Filter: lo0.0-i
```

Name	Bytes	Packets
accept-established-tcp-fetch-lo0.0-i	0	0
accept-established-tcp-ftp-lo0.0-i	1044	12
accept-established-tcp-ftp-data-lo0.0-i	256	4
accept-established-tcp-ftp-data-syn-lo0.0-i	1966	24
accept-established-tcp-ssh-lo0.0-i	0	0
accept-established-tcp-telnet-lo0.0-i	1068	24
accept-established-udp-ephemeral-lo0.0-i	0	0
accept-icmp-lo0.0-i	7072	16
accept-ntp-lo0.0-i	0	0
accept-ntp-server-lo0.0-i	76	1
accept-snmp-lo0.0-i	0	0
accept-ssh-lo0.0-i	8773646	11310
accept-traceroute-icmp-lo0.0-i	276	3
accept-traceroute-tcp-lo0.0-i	0	0
accept-traceroute-udp-lo0.0-i	0	0
accept-web-lo0.0-i	0	0
discard-all-TTL_1-unknown-lo0.0-i	0	0
discard-icmp-lo0.0-i	0	0
discard-netbios-lo0.0-i	2730	35
discard-tcp-lo0.0-i	0	0
discard-udp-lo0.0-i	0	0
discard-unknown-lo0.0-i	0	0
no-icmp-fragments-lo0.0-i	30592	24

Policers:

Name	Packets
management-1m-accept-ntp-lo0.0-i	0
management-1m-accept-ntp-server-lo0.0-i	0
management-1m-accept-traceroute-icmp-lo0.0-i	0
management-1m-accept-traceroute-tcp-lo0.0-i	0
management-1m-accept-traceroute-udp-lo0.0-i	0
management-5m-accept-established-tcp-fetch-lo0.0-i	0
management-5m-accept-established-tcp-ftp-lo0.0-i	0

```

management-5m-accept-established-tcp-ftp-data-100.0-i          0
management-5m-accept-established-tcp-ftp-data-syn-100.0-i    0
management-5m-accept-established-tcp-ssh-100.0-i            0
management-5m-accept-established-tcp-telnet-100.0-i         0
management-5m-accept-established-udp-ephemeral-100.0-i      0
management-5m-accept-icmp-100.0-i                          0
management-5m-accept-snmp-100.0-i                          0
management-5m-accept-ssh-100.0-i                            102
management-5m-accept-web-100.0-i                            0

```

dhanks@MX80>

A lot of things changed now that there's traffic being sent to and from the routing engine. Let's take a look at the changes task by task.

Using Your PC or Another Router, Use the traceroute Command to the Router

This task sent traffic towards the routing engine using the traceroute command and the author's implementation of traceroute used ICMP. You can see that the counter `accept-traceroute-icmp-100.0-i` found three packets. This makes sense as traceroute sends three packets by default:

```
accept-traceroute-icmp-100.0-i          276      3
```

Using Your PC or Another Router, Use the SCP Command and Copy a File to the Router

This task is two-fold: it would increment the counter `accept-ssh-100.0-i`, but also invoke the policer that's installed in the firewall filter `accept-ssh`, which is set to police traffic exceeding a bandwidth limit of 5 Mbps:

```

Counters:
accept-ssh-100.0-i          8773646    11310
Policers:
management-5m-accept-ssh-100.0-i          102

```

Using Your PC or Another Router, Use the Ping Command to the Router

The ping command incremented the counter `accept-icmp-100.0-i`. No surprise there:

```
accept-icmp-100.0-i          7072     16
```

Using Your PC or Another Router, Use the Ping Command, but Change the Packet Size to 9000

This is an interesting task because the ping command is used in such a way that it would fragment the ICMP traffic. The firewall filter `accept-`

icmp and the first term no-icmp-fragments found these packets, incremented the counter, and discarded them:

```
no-icmp-fragments-lo0.0-i          30592      24
```

Use the show firewall log to See the Packet Headers of the Discarded Packets

```
dhanks@MX80> show firewall log
```

Log :

Time	Filter	Action	Interface	Protocol	Src Addr	Dest Addr
03:43:11	lo0.0-i	D	fxp0.0	ICMP	172.16.1.100	172.16.1.11
03:43:11	lo0.0-i	D	fxp0.0	ICMP	172.16.1.100	172.16.1.11
03:43:11	lo0.0-i	D	fxp0.0	ICMP	172.16.1.100	172.16.1.11
03:43:11	lo0.0-i	D	fxp0.0	ICMP	172.16.1.100	172.16.1.11

To get more detailed information, use the show firewall log detail command:

```
dhanks@MX80> show firewall log detail
```

```
Time of Log: 2011-04-03 03:43:11 UTC, Filter: lo0.0-i, Filter action: discard, Name of
interface: fxp0.0
Name of protocol: ICMP, Packet Length: 48288, Source address: 172.16.1.100, Destination
address: 172.16.1.11
ICMP type: 0, ICMP code: 0
Time of Log: 2011-04-03 03:43:11 UTC, Filter: lo0.0-i, Filter action: discard, Name of
interface: fxp0.0
Name of protocol: ICMP, Packet Length: 48288, Source address: 172.16.1.100, Destination
address: 172.16.1.11
ICMP type: 0, ICMP code: 0
Time of Log: 2011-04-03 03:43:11 UTC, Filter: lo0.0-i, Filter action: discard, Name of
interface: fxp0.0
Name of protocol: ICMP, Packet Length: 48288, Source address: 172.16.1.100, Destination
address: 172.16.1.11
ICMP type: 0, ICMP code: 0
Time of Log: 2011-04-03 03:43:11 UTC, Filter: lo0.0-i, Filter action: discard, Name of
interface: fxp0.0
Name of protocol: ICMP, Packet Length: 48288, Source address: 172.16.1.100, Destination
address: 172.16.1.11
ICMP type: 0, ICMP code: 0
```

Using the Router, telnet to Another Router or Device

This task was a bit different, because the router itself was used to generate traffic. The return telnet traffic was caught by the firewall filter accept-established term accept-established-tcp-telnet-lo0.0-i:

```
accept-established-tcp-telnet-lo0.0-i      1068      24
```

Using the Router, FTP to Another Router or Device and Copy a File

Again, traffic was generated from the router itself, but this time using FTP. This is a unique task because the firewall filter `accept-established` has three separate terms to match the return traffic:

- `accept-established-tcp-ftp`: matches the FTP CMD port return traffic.
- `accept-established-tcp-ftp-data-syn`: matches the first SYN packet initiated by the FTP server back to Junos.
- `accept-established-tcp-ftp-data`: matches the return data used by the FTP data port.

```
accept-established-tcp-ftp-100.0-i      1044    12
accept-established-tcp-ftp-data-100.0-i  256     4
accept-established-tcp-ftp-data-syn-100.0-i 1966    24
```

Summary

With this simple framework it's easy to see that the router only accepts traffic that it's interested in. We also gain visibility into the routing engine and understand what type of traffic has been accepted, discarded, and policed.

Advanced Security Policy

Now it's time to introduce more routers and protocols to the mix. We'll create a firewall framework to support the following protocols:

- OSPF
- RIP
- BGP
- BFD
- LDP
- RSVP

This should be more than enough to demonstrate the power and flexibility of the Junos security framework.

For those of you who are following along on your lab router, this book is using a Juniper MX80 router, which supports a router virtualization feature called logical systems. R3 and R4 are logical systems defined within the same physical MX80 as shown in Figure 6.1.

NOTE It's possible to follow along using two physical routers, just make sure that each router has the firewall framework loaded.

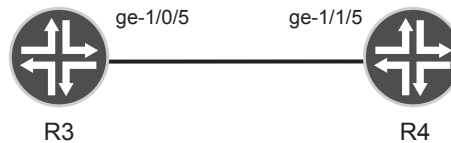


Figure 6.1 Topology Used for the Advanced Security Policy⁹⁹

Quick Look at the Protocol Configuration

R3 and R4 are running quite a number of protocols. Let's take a quick look at how they're configured. You can just look at R3's `protocol` stanza, because R4 has an identical configuration, minus the interface names:

```
[edit]
dhanks@MX80:R3# show protocols
rsvp {
  interface ge-1/0/5.0;
}
mpls {
  label-switched-path R3-to-R4 {
    to 10.0.3.4;
  }
  interface ge-1/0/5.0;
}
bgp {
  export rip-export;
  group 1 {
    type internal;
    local-address 10.0.3.3;
    local-as 65000;
    neighbor 10.0.3.4;
  }
}
ospf {
  area 0.0.0.0 {
    interface ge-1/0/5.0 {
      bfd-liveness-detection {
        minimum-interval 150;
        multiplier 3;
      }
    }
  }
}
```



```

        interface lo0.3 {
            passive;
        }
    }
}
ldp {
    interface lo0.3;
    neighbor 10.0.3.4;
}
rip {
    group 1 {
        export rip-export;
        neighbor ge-1/0/5.0;
    }
}
}

```

Verify the Protocols

Before applying the security policy, you should verify that all of the protocols are configured correctly, adjacencies are up, and things are working as expected.

OSPF

```
dhanks@MX80:R3> show ospf neighbor
```

Address	Interface	State	ID	Pri	Dead
10.0.2.6	ge-1/0/5.0	Full	10.0.3.4	128	37

```
dhanks@MX80:R3>
```

OSPF neighbor is up with the FULL state. Looks good.

RIP

```
dhanks@MX80:R3> show rip neighbor
```

Neighbor	Source State	Source Address	Destination Address	Send Mode	Receive Mode	In Met
ge-1/0/5.0	Up	10.0.2.5	224.0.0.9	mcast	both	1

```
dhanks@MX80:R3>
```

RIP is up and running.

BGP

```
dhanks@MX80:R3> show bgp summary
```

```
Groups: 1 Peers: 1 Down peers: 0
Table      Tot Paths  Act Paths  Suppressed  History  Damp State  Pending
```

```
inet.0          1      0      0      0      0      0
Peer           AS      InPkt  OutPkt  OutQ   Flaps  Last Up/Dwn State|#Active/
Received/Accepted/Damped...
10.0.3.4       65000  1073   1071    0      0      8:01:42 0/1/1/0
0/0/0/0
```

dhanks@MX80:R3>

BGP is up and in the ESTABLISHED state.

BFD

dhanks@MX80:R3> **show bfd session**

```
Address          State      Interface      Detect Time   Transmit Interval Multiplier
10.0.2.6         Up         ge-1/0/5.0     0.450      0.150      3
```

1 sessions, 1 clients

Cumulative transmit rate 6.7 pps, cumulative receive rate 6.7 pps

dhanks@MX80:R3>

BFD is up and transmitting and receiving packets every 150ms.

LDP

dhanks@MX80:R3> **show ldp neighbor**

```
Address          Interface      Label space ID      Hold time
10.0.3.4         lo0.3         10.0.3.4:0         36
```

dhanks@MX80:R3> **show ldp session**

```
Address          State      Connection      Hold time
10.0.3.4         Operational Open            26
```

dhanks@MX80:R3>

You can see the LDP neighbor and the session is open and operational.

RSVP

dhanks@MX80:R3> **show rsvp neighbor**

RSVP neighbor: 1 learned

```
Address          Idle Up/Dn LastChange HelloInt HelloTx/Rx MsgRcvd
10.0.2.6         5 1/0   8:04:38   9 3216/3216 1316
```

dhanks@MX80:R3> °

Ingress RSVP: 1 sessions

```
To          From          State      Rt Style Labelin Labelout LSPname
10.0.3.4    10.0.3.3     Up         0 1 FF      -        3 R3-to-R4
Total 1 displayed, Up 1, Down 0
```

Egress RSVP: 1 sessions

```

To           From           State  Rt Style Labelin Labelout LSPname
10.0.3.3     10.0.3.4       Up     0 1 FF      3      - R4-to-R3
Total 1 displayed, Up 1, Down 0

```

```

Transit RSVP: 0 sessions
Total 0 displayed, Up 0, Down 0

```

```
dhanks@MX80:R3>
```

The RSVP neighbor is up and passing traffic. You can also see that there are ingress and egress sessions that are up.

Apply Security Policy

Now that all the protocols are configured and working as expected, let's move on to applying the firewall filters to the loopbacks of routers R3 and R4.

R3

Although R3 is running a lot of protocols, the firewall filter chain is very simple. We'll use the following firewall filters:

- `accept-common-services`
- `accept-ospf`
- `accept-rip`
- `accept-bfd`
- `accept-bgp`
- `accept-ldp`
- `accept-rsvp`
- `accept-established`
- `discard-all`

Maintaining the firewall order above, apply them to R3's routing engine on interface `lo0.3` like so:

```

[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list accept-common-
services

```

```

[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list accept-ospf

```

```
[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list accept-rip

[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list accept-bfd

[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list accept-bgp

[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list accept-ldp

[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list accept-rsvp

[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list accept-established

[edit]
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list discard-all

[edit]
dhanks@MX80:R3# commit
```

R4

Using the same firewall filters and order, apply them to R4's routing engine on interface lo0.4 like this:

```
[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list accept-common-
services

[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list accept-ospf

[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list accept-rip

[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list accept-bfd

[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list accept-bgp

[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list accept-ldp

[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list accept-rsvp
```

```
[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list accept-established
```

```
[edit]
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list discard-all
```

```
[edit]
dhanks@MX80:R4# commit
```

Verify R3 and R4

Now that R3 and R4 have a security framework applied to the routing engine, let's double check that all the protocols are working as expected again.

```
dhanks@MX80:R3> show ospf neighbor
```

Address	Interface	State	ID	Pri	Dead
10.0.2.6	ge-1/0/5.0	Full	10.0.3.4	128	37

```
dhanks@MX80:R3>
```

```
dhanks@MX80:R3> show rip neighbor
```

Neighbor	State	Source Address	Destination Address	Send Mode	Receive Mode	In Met
ge-1/0/5.0	Up	10.0.2.5	224.0.0.9	mcast	both	1

```
dhanks@MX80:R3>
```

```
dhanks@MX80:R3> show bgp summary
```

```
Groups: 1 Peers: 1 Down peers: 0
Table Tot Paths Act Paths Suppressed History Damp State Pending
inet.0 1 0 0 0 0 0 Last Up/Dwn State|#Active/
Peer AS InPkt OutPkt OutQ Flaps Last Up/Dwn State|#Active/
Received/Accepted/Damped...
10.0.3.4 65000 1125 1124 0 0 8:25:24 0/1/1/0
0/0/0/0
```

```
dhanks@MX80:R3>
```

```
dhanks@MX80:R3> show bfd session
```

Address	State	Interface	Detect Time	Transmit Interval	Multiplier
10.0.2.6	Up	ge-1/0/5.0	0.450	0.150	3

```
1 sessions, 1 clients
```

```
Cumulative transmit rate 6.7 pps, cumulative receive rate 6.7 pps
```

```
dhanks@MX80:R3>
```

```
dhanks@MX80:R3> show ldp neighbor
```

```

Address          Interface      Label space ID      Hold time
10.0.3.4         lo0.3         10.0.3.4:0         36

```

```
dhanks@MX80:R3> show ldp session
```

```

Address          State          Connection          Hold time
10.0.3.4         Operational   Open                26

```

```
dhanks@MX80:R3>
```

```
dhanks@MX80:R3> show rsvp neighbor
```

```
RSVP neighbor: 1 learned
```

```

Address          Idle Up/Dn LastChange HelloInt HelloTx/Rx MsgRcvd
10.0.2.6         0 1/0   8:26:51   9 3363/3363 1378

```

```
dhanks@MX80:R3> show rsvp session
```

```
Ingress RSVP: 1 sessions
```

```

To          From          State   Rt Style Labelin Labelout LSPname
10.0.3.4    10.0.3.3      Up      0 1 FF      -        3 R3-to-R4
Total 1 displayed, Up 1, Down 0

```

```
Egress RSVP: 1 sessions
```

```

To          From          State   Rt Style Labelin Labelout LSPname
10.0.3.3    10.0.3.4      Up      0 1 FF      3        - R4-to-R3
Total 1 displayed, Up 1, Down 0

```

```
Transit RSVP: 0 sessions
```

```
Total 0 displayed, Up 0, Down 0
```

```
dhanks@MX80:R3>
```

Everything is up and running as expected. It looks like the security framework applied to R3 and R4 is working properly.

Viewing Counters, Logs, and Policers

Let's take a look at the counters, logs, and policers on R3. You should expect to see the counters matching protocol traffic being *incremented*.

NOTE The author is using a virtualization feature called logical systems. All `show firewall` commands must be performed from the physical router and must then reference the specific logical system filter. The filter `10.0.3-i` refers to the logical-system R3.

```
dhanks@MX80> show firewall filter 10.0.3-i
```

```

Filter: 10.0.3-i
Counters:

```

Name	Bytes	Packets
accept-bfd-100.3-i	11764012	226231
accept-bgp-100.3-i	139108	2259
accept-established-tcp-fetch-100.3-i	0	0
accept-established-tcp-ftp-100.3-i	0	0
accept-established-tcp-ftp-data-100.3-i	0	0
accept-established-tcp-ftp-data-syn-100.3-i	0	0
accept-established-tcp-ssh-100.3-i	0	0
accept-established-tcp-telnet-100.3-i	0	0
accept-established-udp-ephemeral-100.3-i	0	0
accept-icmp-100.3-i	84168	1002
accept-ldp-discover-100.3-i	0	0
accept-ldp-igmp-100.3-i	0	0
accept-ldp-unicast-100.3-i	357312	5810
accept-ntp-100.3-i	0	0
accept-ntp-server-100.3-i	0	0
accept-ospf-100.3-i	247044	3621
accept-rip-100.3-i	55952	1076
accept-rip-igmp-100.3-i	32	1
accept-rsvp-100.3-i	428488	4764
accept-snmp-100.3-i	0	0
accept-ssh-100.3-i	64	1
accept-tldp-discover-100.3-i	164150	2345
accept-traceroute-icmp-100.3-i	0	0
accept-traceroute-tcp-100.3-i	0	0
accept-traceroute-udp-100.3-i	960	24
accept-web-100.3-i	0	0
discard-all-TTL_1-unknown-100.3-i	0	0
discard-icmp-100.3-i	0	0
discard-netbios-100.3-i	0	0
discard-tcp-100.3-i	0	0
discard-udp-100.3-i	0	0
discard-unknown-100.3-i	0	0
no-icmp-fragments-100.3-i	0	0
Policers:		
Name		Packets
management-1m-accept-ntp-100.3-i		0
management-1m-accept-ntp-server-100.3-i		0
management-1m-accept-traceroute-icmp-100.3-i		0
management-1m-accept-traceroute-tcp-100.3-i		0
management-1m-accept-traceroute-udp-100.3-i		0
management-5m-accept-established-tcp-fetch-100.3-i		0
management-5m-accept-established-tcp-ftp-100.3-i		0
management-5m-accept-established-tcp-ftp-data-100.3-i		0
management-5m-accept-established-tcp-ftp-data-syn-100.3-i		0
management-5m-accept-established-tcp-ssh-100.3-i		0
management-5m-accept-established-tcp-telnet-100.3-i		0
management-5m-accept-established-udp-ephemeral-100.3-i		0
management-5m-accept-icmp-100.3-i		0
management-5m-accept-snmp-100.3-i		0
management-5m-accept-ssh-100.3-i		0

```
management-5m-accept-web-100.3-i 0
```

```
dhanks@MX80>
```

OSPF

The firewall filter `accept-ospf` is matching all of the OSPF traffic and incrementing the counter `accept-ospf-100.3-i`. You can see that the number of packets is incremented for every OSPF Hello packet:

```
dhanks@MX80> show firewall filter 100.3-i counter accept-ospf-100.3-i
```

```
Filter: 100.3-i
Counters:
Name                               Bytes           Packets
accept-ospf-100.3-i                250320          3669
```

Wait a few moments for an OSPF Hello Packet ...

```
dhanks@MX80> show firewall filter 100.3-i counter accept-ospf-100.3-i
```

```
Filter: 100.3-i
Counters:
Name                               Bytes           Packets
accept-ospf-100.3-i                250388          3670
```

```
dhanks@MX80>
```

Notice how the first command shows a total of 3669 packets have been counted. A few seconds later the same command was executed and the packets have increased to 3670.

RIP

RIP is being matched by the firewall filter `accept-rip` and the counter `accept-rip-100.3-i` is being incremented with every RIP update:

```
dhanks@MX80> show firewall filter 100.3-i counter accept-rip-100.3-i
```

```
Filter: 100.3-i
Counters:
Name                               Bytes           Packets
accept-rip-100.3-i                 56784           1092
```

Wait a few moments for a RIP update ...

```
dhanks@MX80> show firewall filter 100.3-i counter accept-rip-100.3-i
```

```
Filter: 100.3-i
Counters:
Name                               Bytes           Packets
accept-rip-100.3-i                 56836           1093
```

```
dhanks@MX80>
```


RIP started out with 1092 packets and after a moment the packets have increased to 1093.

BGP

The firewall filter *accept-bgp* is matching all of the BGP traffic being sent to R3. The counter *accept-bgp-lo0.3-i* is also being incremented with every BGP packet:

```
dhanks@MX80> show firewall filter lo0.3-i counter accept-bgp-lo0.3-i
```

```
Filter: lo0.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-bgp-lo0.3-i	141814	2303

Wait a few moments ...

```
dhanks@MX80> show firewall filter lo0.3-i counter accept-bgp-lo0.3-i
```

```
Filter: lo0.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-bgp-lo0.3-i	141885	2304

```
dhanks@MX80>
```

BFD

BFD is a light-weight protocol, but it sends packets at a high-rate. The firewall filter *accept-bfd* is matching all of the BFD traffic and the counter *accept-bfd-lo0.3-i* is being incremented with every BFD echo:

```
dhanks@MX80> show firewall filter lo0.3-i counter accept-bfd-lo0.3-i
```

```
Filter: lo0.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-bfd-lo0.3-i	12037272	231486

Wait a few moments for some BFD echo packets ...

```
dhanks@MX80> show firewall filter lo0.3-i counter accept-bfd-lo0.3-i
```

```
Filter: lo0.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-bfd-lo0.3-i	12038832	231516

```
dhanks@MX80>
```

LDP

The firewall filter *accept-ldp* has multiple terms to match LDP depending on how it's configured. Here, a targeted LDP is used in the configuration to demonstrate how the discover packets are separated from the unicast traffic. The counter *accept-tldp-discover-100.3-i* is incremented for every targeted LDP discovery packet. The counter *accept-ldp-unicast-100.3-i* is incremented for every LDP session packet:

```
dhanks@MX80> show firewall filter 100.3-i counter accept-ldp-unicast-100.3-i
```

```
Filter: 100.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-ldp-unicast-100.3-i	365720	5944

Wait a few moments for a LDP Hello packet ...

```
dhanks@MX80> show firewall filter 100.3-i counter accept-ldp-unicast-100.3-i
```

```
Filter: 100.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-ldp-unicast-100.3-i	365842	5946

Now let's check out targeted LDP:

```
dhanks@MX80> show firewall filter 100.3-i counter accept-tldp-discover-100.3-i
```

```
Filter: 100.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-tldp-discover-100.3-i	168910	2413

Wait a few moments for an LDP discover packet.

```
dhanks@MX80> show firewall filter 100.3-i counter accept-tldp-discover-100.3-i
```

```
Filter: 100.3-i
```

```
Counters:
```

Name	Bytes	Packets
accept-tldp-discover-100.3-i	168980	2414

```
dhanks@MX80>
```

RSVP

RSVP is being matched by the firewall filter *accept-rsvp*. The counter *accept-rsvp-lo0.3-i* is also being incremented with every packet:

```
dhanks@MX80> show firewall filter lo0.3-i counter accept-rsvp-lo0.3-i
```

```
Filter: lo0.3-i
Counters:
Name                               Bytes           Packets
accept-rsvp-lo0.3-i                443084          4928
```

Wait a few moments for a RSVP hello packet ...

```
dhanks@MX80> show firewall filter lo0.3-i counter accept-rsvp-lo0.3-i
```

```
Filter: lo0.3-i
Counters:
Name                               Bytes           Packets
accept-rsvp-lo0.3-i                443308          4929
```

```
dhanks@MX80>
```

Summary

Although our topology only contained two routers, you could see the number of protocols increased the complexity. Applying the customized firewall filters literally takes only *one command per router*:

```
dhanks@MX80:R3# set interfaces lo0.3 family inet filter input-list [ accept-common-
services accept-ospf accept-rip accept-bfd accept-bgp accept-ldp accept-rsvp accept-
established discard-all ]
```

```
dhanks@MX80:R4# set interfaces lo0.4 family inet filter input-list [ accept-common-
services accept-ospf accept-rip accept-bfd accept-bgp accept-ldp accept-rsvp accept-
established discard-all ]
```

Each router was able to fully function using a wide variety of protocols, while at the same time blocking any traffic that wasn't explicitly configured or that exceeded a policer's bandwidth limit.

No matter how complex the problem, it can always be broken down into simple building blocks. It's much easier to pick and choose firewall filters from a framework than trying to build and maintain large

“protect-re” firewall filters that have to be customized for every router. Maintaining such a large “protect-re” requires too much work for every change in the network, because terms have to be copied into the configuration and moved around to ensure proper evaluation.

As your network changes, modifying the firewall framework is a trivial task. Because all of the building blocks are already preloaded on the router, it’s just a matter of applying them to the loopback interface.

The framework itself can and should be modified to meet the requirements of your network. Try it yourself.

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If you've followed along with this book on a router, you know that creating a framework of firewall filters designed to secure the routing engine is hard work. Each protocol and service has their own little nuances, sometimes making a stateless firewall filter near impossible. It's difficult to gauge how secure versus how broad to make each filter. There's a tradeoff between portability and security. In this book, it was decided to err on the side of caution and chose portability over security in this book's examples. Feel free to tighten the filters up to best secure your own environment.

Lessons Learned

Hopefully this book has showed you that maintaining a simple `filter input-list` is more natural than maintaining a single, gigantic "protect-re" that's over 800 lines long. The simplicity becomes more apparent as this solution is scaled across multiple routers.

Junos configuration automation using `apply-path` feature is a very powerful tool. Each firewall filter used a prefix list leveraging `apply-path` where possible. Using `apply-path` pushes the grunt work of updating changes into prefix lists off the shoulders of the administrator and onto Junos. When administrators perform scheduled maintenance at 3am to add twenty BGP neighbors, and turn up four new interfaces, the last thing the administrator needs to worry about modifying are the firewall filters responsible for securing the routing engine.

You should have also noticed that firewall filters that break down the incoming traffic by protocol, function, and role, provide better analysis and detailed statistics for the administrator to troubleshoot problems. There's no such thing as too much data when it's properly sorted and categorized.

The rest of this appendix contains the framework configuration for your inspection and use. First a few words of caution.

- Care should always be taken when applying firewall filters to the routing engine.
- Always double check your work.
- Use the `show | compare` command liberally, to fully understand what's being changed.
- In production environments use the final firewall filter `accept-all` instead of the `discard-all` command, which allows you to

spend time making sure that all your previous firewall filters in the chain are working as expected.

- The router should be verified to work properly without incrementing counters in the filter `accept-all`. After the verification is complete, it's recommended to immediately remove the filter `accept-all` and replace it with `discard-all` to ensure the routing engine is secure.

But if you've read this book you know all this. Congratulations on securing your Junos routing engine!

TIP

If you visit www.juniper.net/dayone, and then follow the path to this book's download page, you'll find a free Copy and Paste edition of this book. Its rich-text format allows the file to be opened in various text editors for easy copying and pasting of the book's configurations, such as the one that follows in this appendix.

Framework Configuration

```
policy-options {
  prefix-list router-ipv4 {
    apply-path "interfaces <*> unit <*> family inet address <*>";
  }
  prefix-list bgp-neighbors {
    apply-path "protocols bgp group <*> neighbor <*>";
  }
  prefix-list ospf {
    224.0.0.5/32;
    224.0.0.6/32;
  }
  prefix-list rfc1918 {
    10.0.0.0/8;
    172.16.0.0/12;
    192.168.0.0/16;
  }
  prefix-list rip {
    224.0.0.9/32;
  }
  prefix-list vrrp {
    224.0.0.18/32;
  }
  prefix-list multicast-all-routers {
    224.0.0.2/32;
  }
  prefix-list router-ipv4-logical-systems {
    apply-path "logical-systems <*> interfaces <*> unit <*> family inet address
```

```

<*>";
}
prefix-list bgp-neighbors-logical-systems {
    apply-path "logical-systems <*> protocols bgp group <*> neighbor <*>";
}
prefix-list radius-servers {
    apply-path "system radius-server <*>";
}
prefix-list tacas-servers {
    apply-path "system tacplus-server <*>";
}
prefix-list ntp-server {
    apply-path "system ntp server <*>";
}
prefix-list snmp-client-lists {
    apply-path "snmp client-list <*> <*>";
}
prefix-list snmp-community-clients {
    apply-path "snmp community <*> clients <*>";
}
prefix-list localhost {
    127.0.0.1/32;
}
prefix-list ntp-server-peers {
    apply-path "system ntp peer <*>";
}
prefix-list dns-servers {
    apply-path "system name-server <*>";
}
}
firewall {
    family inet {
        prefix-action management-police-set {
            apply-flags omit;
            policer management-1m;
            count;
            filter-specific;
            subnet-prefix-length 24;
            destination-prefix-length 32;
        }
        prefix-action management-high-police-set {
            apply-flags omit;
            policer management-5m;
            count;
            filter-specific;
            subnet-prefix-length 24;
            destination-prefix-length 32;
        }
        filter accept-bgp {
            apply-flags omit;
            term accept-bgp {
                from {

```



```
        source-prefix-list {
            bgp-neighbors;
            bgp-neighbors-logical-systems;
        }
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        protocol tcp;
        port bgp;
    }
    then {
        count accept-bgp;
        accept;
    }
}
}
filter accept-ospf {
    apply-flags omit;
    term accept-ospf {
        from {
            source-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            destination-prefix-list {
                router-ipv4;
                ospf;
                router-ipv4-logical-systems;
            }
        }
        protocol ospf;
    }
    then {
        count accept-ospf;
        accept;
    }
}
}
filter accept-rip {
    apply-flags omit;
    term accept-rip {
        from {
            source-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            destination-prefix-list {
                rip;
            }
        }
        protocol udp;
        destination-port rip;
    }
}
```

```
    }
    then {
        count accept-rip;
        accept;
    }
}
term accept-rip-igmp {
    from {
        source-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        destination-prefix-list {
            rip;
        }
        protocol igmp;
    }
    then {
        count accept-rip-igmp;
        accept;
    }
}
}
filter accept-vrrp {
    apply-flags omit;
    term accept-vrrp {
        from {
            source-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            destination-prefix-list {
                vrrp;
            }
            protocol [ vrrp ah ];
        }
        then {
            count accept-vrrp;
            accept;
        }
    }
}
filter accept-icmp {
    apply-flags omit;
    term no-icmp-fragments {
        from {
            is-fragment;
            protocol icmp;
        }
        then {
            count no-icmp-fragments;
        }
    }
}
```

```

        log;
        discard;
    }
}
term accept-icmp {
    from {
        protocol icmp;
        ttl-except 1;
        icmp-type [ echo-reply echo-request time-exceeded unreachable source-
quench router-advertisement parameter-problem ];
    }
    then {
        policer management-5m;
        count accept-icmp;
        accept;
    }
}
}
filter accept-ssh {
    apply-flags omit;
    term accept-ssh {
        from {
            source-prefix-list {
                rfc1918;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol tcp;
            destination-port ssh;
        }
        then {
            policer management-5m;
            count accept-ssh;
            accept;
        }
    }
}
}
filter accept-snmp {
    apply-flags omit;
    term accept-snmp {
        from {
            source-prefix-list {
                snmp-client-lists;
                snmp-community-clients;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
        }
    }
}
}

```

```
        protocol udp;
        destination-port snmp;
    }
    then {
        policer management-5m;
        count accept-snm;
        accept;
    }
}
filter accept-ntp {
    apply-flags omit;
    term accept-ntp {
        from {
            source-prefix-list {
                ntp-server;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol udp;
            port ntp;
        }
        then {
            policer management-1m;
            count accept-ntp;
            accept;
        }
    }
}
term accept-ntp-peer {
    from {
        source-prefix-list {
            ntp-server-peers;
        }
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        protocol udp;
        destination-port ntp;
    }
    then {
        policer management-1m;
        count accept-ntp-peer;
        accept;
    }
}
term accept-ntp-server {
    from {
        source-prefix-list {
```

```
        rfc1918;
    }
    destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
    }
    protocol udp;
    destination-port ntp;
}
then {
    policer management-1m;
    count accept-ntp-server;
    accept;
}
}
}
filter accept-web {
    apply-flags omit;
    term accept-web {
        from {
            source-prefix-list {
                rfc1918;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol tcp;
            destination-port [ http https ];
        }
        then {
            policer management-5m;
            count accept-web;
            accept;
        }
    }
}
filter discard-all {
    apply-flags omit;
    term discard-ip-options {
        from {
            ip-options any;
        }
        then {
            count discard-ip-options;
            log;
            syslog;
            discard;
        }
    }
}
term discard-TTL_1-unknown {
```

```
    from {
        ttl 1;
    }
    then {
        count discard-all-TTL_1-unknown;
        log;
        syslog;
        discard;
    }
}
term discard-tcp {
    from {
        protocol tcp;
    }
    then {
        count discard-tcp;
        log;
        syslog;
        discard;
    }
}
term discard-netbios {
    from {
        protocol udp;
        destination-port 137;
    }
    then {
        count discard-netbios;
        log;
        syslog;
        discard;
    }
}
term discard-udp {
    from {
        protocol udp;
    }
    then {
        count discard-udp;
        log;
        syslog;
        discard;
    }
}
term discard-icmp {
    from {
        protocol icmp;
    }
    then {
        count discard-icmp;
        log;
    }
}
```

```
        syslog;
        discard;
    }
}
term discard-unknown {
    then {
        count discard-unknown;
        log;
        syslog;
        discard;
    }
}
}
filter accept-traceroute {
    apply-flags omit;
    term accept-traceroute-udp {
        from {
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol udp;
            ttl 1;
            destination-port 33435-33450;
        }
        then {
            policer management-1m;
            count accept-traceroute-udp;
            accept;
        }
    }
    term accept-traceroute-icmp {
        from {
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol icmp;
            ttl 1;
            icmp-type [ echo-request timestamp time-exceeded ];
        }
        then {
            policer management-1m;
            count accept-traceroute-icmp;
            accept;
        }
    }
}
term accept-traceroute-tcp {
    from {
        destination-prefix-list {
            router-ipv4;
```

```
        router-ipv4-logical-systems;
    }
    protocol tcp;
    ttl 1;
}
then {
    policer management-1m;
    count accept-traceroute-tcp;
    accept;
}
}
}
filter accept-igmp {
    apply-flags omit;
    term accept-ospf {
        filter accept-ospf;
    }
    term accept-rip {
        filter accept-rip;
    }
}
filter accept-common-services {
    apply-flags omit;
    term accept-icmp {
        filter accept-icmp;
    }
    term accept-traceroute {
        filter accept-traceroute;
    }
    term accept-ssh {
        filter accept-ssh;
    }
    term accept-snmp {
        filter accept-snmp;
    }
    term accept-ntp {
        filter accept-ntp;
    }
    term accept-web {
        filter accept-web;
    }
    term accept-dns {
        filter accept-dns;
    }
}
filter accept-bfd {
    apply-flags omit;
    term accept-bfd {
```



```
    from {
      source-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      protocol udp;
      source-port 49152-65535;
      destination-port 3784-3785;
    }
    then {
      count accept-bfd;
      accept;
    }
  }
}
filter accept-ldp {
  apply-flags omit;
  term accept-ldp-discover {
    from {
      source-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      destination-prefix-list {
        multicast-all-routers;
      }
      protocol udp;
      destination-port ldp;
    }
    then {
      count accept-ldp-discover;
      accept;
    }
  }
}
term accept-ldp-unicast {
  from {
    source-prefix-list {
      router-ipv4;
      router-ipv4-logical-systems;
    }
    destination-prefix-list {
      router-ipv4;
      router-ipv4-logical-systems;
    }
  }
  protocol tcp;
}
```

```
        port ldp;
    }
    then {
        count accept-ldp-unicast;
        accept;
    }
}
term accept-tldp-discover {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        protocol udp;
        destination-port ldp;
    }
    then {
        count accept-tldp-discover;
        accept;
    }
}
term accept-ldp-igmp {
    from {
        source-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        destination-prefix-list {
            multicast-all-routers;
        }
        protocol igmp;
    }
    then {
        count accept-ldp-igmp;
        accept;
    }
}
}
filter accept-ftp {
    apply-flags omit;
    term accept-ftp {
        from {
            source-prefix-list {
                rfc1918;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
        }
    }
}
```

```
        protocol tcp;
        port [ ftp ftp-data ];
    }
    then {
        policer management-5m;
        count accept-ftp;
        accept;
    }
}
filter accept-rsvp {
    apply-flags omit;
    term accept-rsvp {
        from {
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol rsvp;
        }
        then {
            count accept-rsvp;
            accept;
        }
    }
}
filter accept-radius {
    apply-flags omit;
    term accept-radius {
        from {
            source-prefix-list {
                radius-servers;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol udp;
            source-port [ radacct radius ];
            tcp-established;
        }
        then {
            policer management-1m;
            count accept-radius;
            accept;
        }
    }
}
filter accept-tacas {
    apply-flags omit;
    term accept-tacas {
```

```
        from {
            source-prefix-list {
                tacas-servers;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol [ tcp udp ];
            source-port [ tacacs tacacs-ds ];
            tcp-established;
        }
        then {
            policer management-1m;
            count accept-tacas;
            accept;
        }
    }
}
filter accept-remote-auth {
    apply-flags omit;
    term accept-radius {
        filter accept-radius;
    }
    term accept-tacas {
        filter accept-tacas;
    }
}
filter accept-telnet {
    apply-flags omit;
    term accept-telnet {
        from {
            source-prefix-list {
                rfc1918;
            }
            destination-prefix-list {
                router-ipv4;
                router-ipv4-logical-systems;
            }
            protocol tcp;
            destination-port telnet;
        }
        then {
            policer management-1m;
            count accept-telnet;
            accept;
        }
    }
}
filter accept-dns {
    apply-flags omit;
```

```
term accept-dns {
  from {
    source-prefix-list {
      dns-servers;
    }
    destination-prefix-list {
      router-ipv4;
      router-ipv4-logical-systems;
    }
    protocol udp;
    source-port 53;
  }
  then {
    policer management-1m;
    count accept-dns;
    accept;
  }
}
filter accept-ldp-rsvp {
  apply-flags omit;
  term accept-ldp {
    filter accept-ldp;
  }
  term accept-rsvp {
    filter accept-rsvp;
  }
}
filter accept-established {
  apply-flags omit;
  term accept-established-tcp-ssh {
    from {
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      source-port ssh;
      tcp-established;
    }
    then {
      policer management-5m;
      count accept-established-tcp-ssh;
      accept;
    }
  }
}
term accept-established-tcp-ftp {
  from {
    destination-prefix-list {
      router-ipv4;
      router-ipv4-logical-systems;
    }
  }
  source-port ftp;
}
```

```
        tcp-established;
    }
    then {
        policer management-5m;
        count accept-established-tcp-ftp;
        accept;
    }
}
term accept-established-tcp-ftp-data-syn {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        source-port ftp-data;
        tcp-initial;
    }
    then {
        policer management-5m;
        count accept-established-tcp-ftp-data-syn;
        accept;
    }
}
term accept-established-tcp-ftp-data {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        source-port ftp-data;
        tcp-established;
    }
    then {
        policer management-5m;
        count accept-established-tcp-ftp-data;
        accept;
    }
}
term accept-established-tcp-telnet {
    from {
        destination-prefix-list {
            router-ipv4;
            router-ipv4-logical-systems;
        }
        source-port telnet;
        tcp-established;
    }
    then {
        policer management-5m;
        count accept-established-tcp-telnet;
        accept;
    }
}
```

```
    }
  }
  term accept-established-tcp-fetch {
    from {
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      source-port [ http https ];
      tcp-established;
    }
    then {
      policer management-5m;
      count accept-established-tcp-fetch;
      accept;
    }
  }
  term accept-established-udp-ephemeral {
    from {
      destination-prefix-list {
        router-ipv4;
        router-ipv4-logical-systems;
      }
      protocol udp;
      destination-port 49152-65535;
    }
    then {
      policer management-5m;
      count accept-established-udp-ephemeral;
      accept;
    }
  }
}
filter accept-all {
  apply-flags omit;
  term accept-all-tcp {
    from {
      protocol tcp;
    }
    then {
      count accept-all-tcp;
      log;
      syslog;
      accept;
    }
  }
  term accept-all-udp {
    from {
      protocol udp;
    }
    then {
      count accept-all-udp;
    }
  }
}
```



```

    then discard;
  }
}

```

R3 Configuration

```

logical-systems {
  R3 {
    interfaces {
      ge-1/0/5 {
        unit 0 {
          family inet {
            address 10.0.2.5/30;
          }
          family iso;
          family mpls;
        }
      }
      lo0 {
        unit 3 {
          family inet {
            filter {
              input-list [ accept-common-services accept-ospf accept-rip
accept-bfd accept-bgp accept-ldp accept-rsvp discard-all ];
            }
            address 10.0.3.3/32;
          }
          family iso {
            address 49.0002.0100.0000.3003.00;
          }
        }
      }
    }
  }
}
protocols {
  rsvp {
    interface ge-1/0/5.0;
    interface all;
  }
  mpls {
    label-switched-path R3-to-R4 {
      to 10.0.3.4;
    }
    interface ge-1/0/5.0;
  }
  bgp {
    export rip-export;
    group 1 {
      type internal;
      local-address 10.0.3.3;
    }
  }
}

```

```
        local-as 65000;
        neighbor 10.0.3.4;
    }
}
isis {
    export isis-export;
    reference-bandwidth 1g;
    lsp-lifetime 3600;
    traffic-engineering {
        family inet {
            shortcuts;
        }
    }
    level 2 {
        authentication-key "$9$vBy8xDjq.5F"; ## SECRET-DATA
        authentication-type simple;
    }
    level 1 wide-metrics-only;
    interface all {
        level 1 disable;
        level 2 {
            hello-authentication-key "$9$0y1C1EydVY2aU"; ## SECRET-DATA
            hello-authentication-type md5;
        }
    }
}
ospf {
    area 0.0.0.0 {
        interface ge-1/0/5.0 {
            bfd-liveness-detection {
                minimum-interval 150;
                multiplier 3;
            }
        }
        interface lo0.3 {
            passive;
        }
    }
}
ldp {
    interface lo0.3;
    neighbor 10.0.3.4;
}
rip {
    group 1 {
        export rip-export;
        neighbor ge-1/0/5.0;
    }
}
}
```

```

policy-options {
  policy-statement isis-export {
    term 1 {
      from {
        route-filter 10.0.5.0/24 exact;
        route-filter 10.0.4.0/22 exact;
      }
      to level 2;
      then accept;
    }
    term 2 {
      from {
        route-filter 10.0.4.0/22 longer;
      }
      to level 2;
      then reject;
    }
  }
  policy-statement rip-export {
    term 1 {
      from protocol aggregate;
      then accept;
    }
  }
}
routing-options {
  static {
    route 5.1.1.1/32 receive;
  }
  aggregate {
    route 10.0.4.0/22;
    route 5.0.0.0/8;
  }
}
}

```

R4 Configuration

```

logical-systems {
  R4 {
    interfaces {
      ge-1/1/5 {
        unit 0 {
          family inet {
            address 10.0.2.6/30;
          }
          family iso;
        }
      }
    }
  }
}

```

```
        family mpls;
    }
}
lo0 {
    unit 4 {
        family inet {
            filter {
                input-list [ accept-common-services accept-ospf accept-rip
accept-bfd accept-bgp accept-ldp accept-rsvp discard-all ];
            }
            address 10.0.3.4/32;
        }
        family iso {
            address 49.0002.0100.0000.3004.00;
        }
    }
}
}
protocols {
    rsvp {
        interface ge-1/1/5.0;
        interface all;
    }
    mpls {
        label-switched-path R4-to-R3 {
            to 10.0.3.3;
        }
        interface ge-1/1/5.0;
    }
    bgp {
        export ee;
        group 1 {
            type internal;
            local-address 10.0.3.4;
            local-as 65000;
            neighbor 10.0.3.3;
        }
    }
    isis {
        export isis-export;
        reference-bandwidth 1g;
        lsp-lifetime 3600;
        traffic-engineering {
            family inet {
                shortcuts;
            }
        }
        level 2 {
            authentication-key "$9$oEZDk9Cu0Ic"; ## SECRET-DATA
            authentication-type simple;
        }
    }
}
```

```
    level 1 wide-metrics-only;
  interface all {
    level 1 disable;
    level 2 {
      hello-authentication-key "$9$hT7yewg4ZGi."; ## SECRET-DATA
      hello-authentication-type md5;
    }
  }
}
ospf {
  export ee;
  area 0.0.0.0 {
    interface ge-1/1/5.0 {
      bfd-liveness-detection {
        minimum-interval 150;
        multiplier 3;
      }
    }
  }
}
ldp {
  interface lo0.4;
  neighbor 10.0.3.3;
}
rip {
  group 1 {
    export ee;
    neighbor ge-1/1/5.0;
  }
}
}
policy-options {
  policy-statement ee {
    term 1 {
      from protocol aggregate;
      then accept;
    }
  }
  policy-statement isis-export {
    term 1 {
      from {
        route-filter 10.0.5.0/24 exact;
        route-filter 10.0.4.0/22 exact;
      }
      to level 2;
      then accept;
    }
    term 2 {
      from {
        route-filter 10.0.4.0/22 longer;
      }
    }
  }
}
```

```
        to level 2;
        then reject;
    }
}
routing-options {
  static {
    route 2.2.2.2/32 receive;
  }
  aggregate {
    route 10.0.4.0/22;
    route 2.0.0.0/8;
  }
}
}
```